

Quasi-geodesics in $\text{Out}(\mathbb{F}_n)$ and their shadows in sub-factors

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ABSTRACT

We study the behaviour of quasi-geodesics in $\text{Out}(\mathbb{F}_n)$. Given an element ϕ in $\text{Out}(\mathbb{F}_n)$, there are several natural paths connecting the origin to ϕ in $\text{Out}(\mathbb{F}_n)$; for example, paths given by Stallings’ folding algorithm and paths induced by the shadow of greedy folding paths in Outer Space. We show that none of these paths is, in general, a quasi-geodesic in $\text{Out}(\mathbb{F}_n)$. In fact, in contrast with the mapping class group setting, we construct examples where any quasi-geodesic in $\text{Out}(\mathbb{F}_n)$ connecting ϕ to the origin will have to backtrack in some free factor of \mathbb{F}_n .

1. Introduction

Let $\mathbb{F}_n = \langle s_1, s_2, \dots, s_n \rangle$ denote the free group of rank n and let $\text{Out}(\mathbb{F}_n)$ denote the group of outer automorphisms of \mathbb{F}_n ,

$$\text{Out}(\mathbb{F}_n) := \text{Aut}(\mathbb{F}_n) / \text{Inn}(\mathbb{F}_n).$$

This group is finitely presented [28]. For example, it can be generated by the set of right transvections and left transvections:

$$\begin{cases} s_i \rightarrow s_i s_k \\ s_j \rightarrow s_j \end{cases} \quad \text{for all } j \neq i \qquad \begin{cases} s_i \rightarrow s_k s_i \\ s_j \rightarrow s_j \end{cases} \quad \text{for all } j \neq i$$

together with elements permuting the basis and elements sending some s_i to its inverse \bar{s}_i . For an explicit presentation see [27]. Equip $\text{Out}(\mathbb{F}_n)$ with the word metric associated to this generating set. We aim to understand the geometry $\text{Out}(\mathbb{F}_n)$ as a metric space. Specifically, we want to understand the quasi-geodesics in $\text{Out}(\mathbb{F}_n)$.

A common way to generate a path connecting a point in $\text{Out}(\mathbb{F}_n)$ to the identity is to use the seminal idea of Stallings’ folding algorithm [30] which we discuss here briefly. One can create a model for $\text{Out}(\mathbb{F}_n)$ by considering the space of graphs x of rank n where the oriented edges are labelled by elements of \mathbb{F}_n inducing an isomorphism $\mu: \pi_1(x) \rightarrow \mathbb{F}_n$ (defined up to conjugation and graph automorphism). We refer to μ as the *marking map*. For another such marked graph x' with marking map $\mu': \pi_1(x') \rightarrow \mathbb{F}_n$, we say x' is obtained from x by a *fold* if there is a quotient map from x to x' ,

$$\text{fold}: x \rightarrow x',$$

inducing an isomorphism $\text{fold}^*: \pi_1(x) \rightarrow \pi_1(x')$ so that $\mu = \mu' \circ \text{fold}^*$.

Let R_0 be a rose with labels s_1, \dots, s_n inducing an isomorphism

$$\mu_0: \pi_1(R_0) \rightarrow \mathbb{F}_n,$$

and, for $\phi \in \text{Out}(\mathbb{F}_n)$, let $R = \phi(R_0)$ be the rose where the marking map

$$\mu: \pi_1(R) \rightarrow \mathbb{F}_n \quad \text{is given by} \quad \mu = \phi \mu_0.$$

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Then, a sequence

$$R = x_m \rightarrow x_{m-1} \cdots \rightarrow x_1 \rightarrow x_0 = R_0$$

of folds produces a path in $\text{Out}(\mathbb{F}_n)$ connecting the identity to ϕ as follows: for $0 < i \leq m$, consider a quotient map $x_i \rightarrow R_i$ by collapsing a spanning sub-tree, resulting in a rose R_i . Also, let $q_i: x_i \rightarrow x_0 = R_0$ be the composition of the folding maps. Then, we have the following diagram of homotopy equivalences:

$$R_0 \rightarrow R_i \leftarrow x_i \rightarrow R_0,$$

where the first arrow is any graph automorphism between the two roses. We associate x_i to the induced map $\phi_i: \pi_1(R_0) \rightarrow \pi_1(R_0)$. The map $\phi_i \in \text{Out}(\mathbb{F}_n)$ is coarsely well defined, depending on the chosen quotient maps $x_i \rightarrow R_i$ and the graph automorphism $R_0 \rightarrow R_i$. But the set of all possible resulting maps has a uniformly bounded diameter in $\text{Out}(\mathbb{F}_n)$ and the distance between ϕ_i and ϕ_{i+1} in $\text{Out}(\mathbb{F}_n)$ is uniformly bounded. That is, we have a coarse path

$$\phi = \phi_m, \dots, \phi_0 = \text{id}$$

in $\text{Out}(\mathbb{F}_n)$ connecting ϕ to the identity which we refer to as a *folding path* in $\text{Out}(\mathbb{F}_n)$.

Stallings folding theorem provides an algorithm for finding a sequence of folds connecting any marked graph x to R_0 where, in each fold, edges are identified according to their labels (see Section 2.4). That is to say, the folds given by the Stallings algorithm are directed by the marking map. A path consisting of a sequence of marking-directed folds is called a *Stallings folding path*. The ideas in Stallings' folding algorithm have found many interesting applications in the study of automorphisms of free groups (see, for example, [12, 14, 16, 17, 34]).

Among the complexes frequently used to study $\text{Out}(\mathbb{F}_n)$ (see, for example [5, 19, 20]) are the free splitting complex and the free factor complex (see Section 2.2 for definitions). Let $\mathcal{S}(\mathcal{A})$ denote the free splitting complex of \mathbb{F}_n . Stallings folding paths can be thought of as paths in $\mathcal{S}(\mathcal{A})$. Handel and Mosher showed that these paths are quasi-geodesics in $\mathcal{S}(\mathcal{A})$ [19]. Let $\mathcal{F}(\mathcal{A})$ denote the complex of free factors of \mathbb{F}_n . Kapovich and Rafi [20] proved that a projection from $\mathcal{S}(\mathcal{A})$ to $\mathcal{F}(\mathcal{A})$ takes geodesics to quasi-geodesics. Also, generalization of Stallings folding paths give rise to geodesics in CV_n , for example, see [15, 29].

However, Stallings folding paths do not in general give efficient paths in $\text{Out}(\mathbb{F}_n)$.

THEOREM A. *For any $K, C > 0$, there exists an element $\phi \in \text{Out}(\mathbb{F}_n)$ such that any Stallings folding path connecting $R = \phi(R_0)$ to R_0 does not yield a (K, C) -quasi-geodesic path in $\text{Out}(\mathbb{F}_n)$.*

Proof. Consider the following automorphism:

$$\phi: \begin{cases} a \rightarrow a \\ b \rightarrow b \\ c \rightarrow c(ab^s)^t. \end{cases}$$

That is, if R_0 is the rose with edge labels a, b and c , then $R = \phi(R_0)$ has labels a, b , and $(ab^s)^t$. Following Stallings folding algorithm, to go from R to R_0 we need to fold the third edge around the edges labelled a and b . At each step, there is only one fold possible. The first and second edges remain unchanged since they are also present in R_0 . This takes $t(s + 1)$ steps. That is, the associated path in $\text{Out}(\mathbb{F}_n)$ has a length comparable to ts .

However, one can see that there is another path connecting R to R_0 that takes $2s + t$ steps, namely

$$\langle a, b, c(ab^s)^t \rangle \longrightarrow \langle ab^s, b, c(ab^s)^t \rangle \longrightarrow \langle ab^s, b, c \rangle \longrightarrow \langle a, b, c \rangle.$$

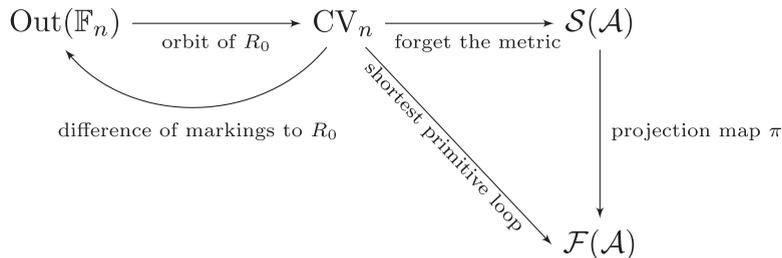


FIGURE 1. A path in either $\text{Out}(\mathbb{F}_n)$ or CV_n has an associated shadow path to other spaces via the maps described here.

Choosing s, t sufficiently large compared with the given K and C , we have shown that the path given by marking-directed folds was not a (K, C) -quasi-geodesic. \square

One can also represent an element of $\text{Out}(\mathbb{F}_n)$ using *train-track maps* (see [6]) and consider a folding sequence according to the train-track structure. Or similarly, consider a geodesic in Culler–Vogtmann Outer Space and take the shadow (that is, the difference of markings to R_0 , see the discussion at the beginning of Section 5 and Figure 1) of it to $\text{Out}(\mathbb{F}_n)$. The same example above shows that none of these paths would, in general, produce a quasi-geodesic in $\text{Out}(\mathbb{F}_n)$.

On the other hand, we observe in the above example, that even along the shorter paths, it takes at least s steps to form or to eliminate an s th power of b and t steps to eliminate a t th power of (ab^s) . We examine this phenomenon through the language of *relative twisting number* [8].

The *relative twisting number* [8] of two labelled graphs x and x' around a loop α measures the difference between x and x' from the point of view of α and is denoted by $\text{tw}_\alpha(x, x')$ (see Definition 3.1). We show that the length of any path in $\text{Out}(\mathbb{F}_n)$ connecting ϕ to the identity is bounded below by the relative twisting number $\text{tw}_\alpha(R, R_0)$, where $R = \phi(R_0)$.

Again, instead of considering a path in the Cayley graph $\text{Out}(\mathbb{F}_n)$, we consider a folding sequence R_m, \dots, R_0 . This is a sequence of labelled graphs where R_i is obtained from R_{i+1} by a fold that it is not necessary coming from Stallings’ algorithm or any train-track structure (see Section 2.4). We also show that if the length of the loop α remains long along the path, it takes longer to twist around α .

THEOREM B. *For any folding sequence R_m, \dots, R_0 and any loop α , we have*

$$m \geq \text{tw}_\alpha(R_0, R_n).$$

Further, if $\ell_{R_i}(\alpha) \geq L > 50$ for every i , then

$$m \geq \text{tw}_\alpha(R_0, R_n) \left(\log_5 \frac{L}{50} \right).$$

One might suspect that, in the above theorem, $\log_5 L$ can be replaced with L . However, we will show that the above inequality is sharp with an example (see Example 3.8).

Theorem B can be viewed in the context of an attempt to have a distance formula for the word length of an element in $\text{Out}(\mathbb{F}_n)$ in analogy with the work of Masur–Minsky in the setting of the mapping class group [26]. Let $S = S_{g,s}$ be a surface of genus g with s punctures and $\text{Mod}(S)$ denote the mapping class group of S , that is, the group of orientation preserving self-homeomorphisms of S up to isotopy. One can try to understand an element $f \in \text{Mod}(S)$ inductively by measuring the contribution of every subsurface to the complexity of f . This is done explicitly as follows: a marking μ_0 on a surface S is a set of simple closed curves that

fill the surface, that is to say, every other curve on S intersects some curve in μ_0 . Masur and Minsky introduced a measure of complexity $d_Y(\mu_0, f(\mu_0))$ between μ_0 and $f(\mu_0)$ called the subsurface projection distance. Namely, they defined a projection map

$$\pi_Y: C(S) \rightarrow C(Y)$$

from the curve graph of S to the curve graph of a sub-surface Y and defined $d_Y(\mu_0, f(\mu_0))$ to be the distance in $C(Y)$ between the projection $\pi_Y(\mu_0)$ and $\pi_Y(f(\mu_0))$ to $C(Y)$. They showed the sum of these subsurface projections is a good estimate for the word length of f (see [26] for more details).

To produce the upper bound for the distance formula, Masur and Minsky constructed a class of quasi-geodesics called *hierarchy paths*, whose lengths is the coarse sum of all subsurface projection distances. An important characteristic of these quasi-geodesics is that they do not backtrack in any subsurface Y . That is, there is a quasi-geodesic,

$$[0, m] \rightarrow \text{Mod}(S), \quad i \rightarrow f_i,$$

in $\text{Mod}(S)$ so that the projection to the curve graph of Y

$$[0, m] \rightarrow C(Y), \quad i \rightarrow \pi_Y(f_i(\mu_0))$$

is a quasi-geodesic for every subsurface Y of S .

There are several analogues for the curve graph in the setting $\text{Out}(\mathbb{F}_n)$, most importantly, the free splitting graph $\mathcal{S}(\mathbb{F}_n)$ and the free factor graph $\mathcal{F}(\mathbb{F}_n)$ both have been shown to be Gromov hyperbolic spaces [5, 19]. For every sub-factor \mathcal{A} of \mathbb{F}_n , we have projection maps [5]

$$\text{Out}(\mathbb{F}_n) \rightarrow \mathcal{S}(\mathcal{A}) \rightarrow \mathcal{F}(\mathcal{A})$$

and it is known that every quasi-geodesic in $\mathcal{S}(\mathbb{F}_n)$ projects to a quasi-geodesic in $\mathcal{F}(\mathbb{F}_n)$ [20]. One may hope to construct quasi-geodesic in $\text{Out}(\mathbb{F}_n)$ where the projections the free splitting or the free factor graph of a sub-factor is always a quasi-geodesic. However, we use Theorem B to prove:

THEOREM C (Quasi-geodesics backtrack in sub-factors). *For given constants $K_1, C_1, K_2, C_2 > 0$ there exists an automorphism, $\phi \in \text{Out}(\mathbb{F}_n)$ such that, for any (K_1, C_1) -quasi-geodesic $p: [0, m] \rightarrow \text{Out}(\mathbb{F}_n)$ with $p(0) = \text{id}$ and $p(m) = \phi$, the shadow $\Theta_{\mathcal{A}} \circ p$ of p in $\mathcal{F}(\mathcal{A})$ is not a (K_2, C_2) -reparameterized quasi-geodesic.*

That is to say, there is $\phi \in \text{Out}(\mathbb{F}_n)$ and a free factor \mathcal{A} , such that every quasi-geodesic connecting ϕ to the identity backtracks in $\mathcal{F}(\mathcal{A})$. In other words, there does not exist a quasi-geodesics between the identity and ϕ that projects to a quasi-geodesics in $\mathcal{F}(\mathcal{A})$. Same is true for $\mathcal{S}(\mathcal{A})$.

Another application of Theorem B is in the understanding of relationships between Outer Space geodesics and $\text{Out}(\mathbb{F}_n)$ geodesics. The Outer Space, denoted CV_n , is the set of metric graphs whose fundamental group is identified with \mathbb{F}_n . It is a CW -complex with $\text{Out}(\mathbb{F}_n)$ action and it was defined by Culler and Vogtmann to study $\text{Out}(\mathbb{F}_n)$ as an analogue of Teichmüller space which has $\text{Mod}(S)$ action [11]. One can project a path in CV_n to a path in $\text{Out}(\mathbb{F}_n)$, by considering the associated difference of markings maps along a path in CV_n . Bestvina and Feighn [5] showed that greedy folding paths in CV_n projects to quasi-geodesics in free factor graphs for all sub-factors. This was used to produce a weak version of a distance formula which give a lower bound for the word length in terms of projection distance to $\mathcal{S}(\mathcal{A})$ [3].

However, it follows from Theorem C that shadows of greedy folding paths are not quasi-geodesics in $\text{Out}(\mathbb{F}_n)$.

THEOREM D. *The shadow of geodesics in CV_n do not in general behave well in $\text{Out}(\mathbb{F}_n)$. More specifically,*

- (i) *for given constants K and C , there are points $x, y \in \text{CV}_n$ such that for every geodesic $[x, y]_{\text{CV}_n}$ in CV_n connecting x to y , its image in $\text{Out}(\mathbb{F}_n)$ is not (K, C) -quasi-geodesic;*
- (ii) *there are points x and y in the thick part of CV_n , such that they are connected by a greedy folding path whose shadow in $\text{Out}(\mathbb{F}_n)$ is not a quasi-geodesic.*

Our methods do not say anything about the projection of quasi-geodesics in $\text{Out}(\mathbb{F}_n)$ to $\mathcal{S}(\mathbb{F}_n)$ or $\mathcal{F}(\mathbb{F}_n)$. It is interesting to know if a geodesic in $\mathcal{S}(\mathbb{F}_n)$ can be used as a guide to construct efficient paths in $\text{Out}(\mathbb{F}_n)$.

QUESTION E. For a given $\phi \in \text{Out}(\mathbb{F}_n)$, does there always exist a quasi-geodesic in $\text{Out}(\mathbb{F}_n)$ connecting ϕ to the identity whose projection to $\mathcal{S}(\mathbb{F}_n)$ is also a quasi-geodesic?

2. Background

2.1. Labelled graphs

Recall from the introduction that a labelled graph x induces an isomorphism $\mu: \pi_1(x) \rightarrow \mathbb{F}_n$ called a *marking*. Two labelled graphs x and x' are equivalent if there is a graph automorphism $f: x \rightarrow x'$ such that the following diagram commutes up to free homotopy:

$$\begin{array}{ccc}
 \pi_1(x') & \xrightarrow{\mu'} & \mathbb{F}_n \\
 f^* \uparrow & \nearrow \mu & \\
 \pi_1(x) & &
 \end{array}$$

Let w denote an element of \mathbb{F}_n . We refer to a conjugacy class $[w]$ of w as a *loop*. For any labelled graph x , and any loop α , there is an immersion of a circle in x representing α which (abusing the notation) we also denote by α . We always assume this immersion to be the shortest in its free homotopy class in terms of the number of edges. The number of edges of a loop α in the given marked graph x is denoted $\ell_x(\alpha)$, and is called the *combinatorial length* of α in x .

Sometimes it is more convenient to work with the universal cover of a marked graph. The universal cover of x is an \mathbb{F}_n -tree (A simplicial tree with free \mathbb{F}_n action). Given such a tree T , an element $w \in \mathbb{F}_n$ in the conjugacy class α acts hyperbolically on T , and we use $\text{axis}_T(w)$ to denote the axis of its action. Consistent with the definition of combinatorial lengths in the graphs, we use $\ell_T(w)$ to denote the number of edges in a fundamental domain of the action of w . Note that this is independent of the choice of $w \in \alpha$. Hence, we can also use the notation $\ell_T(\alpha)$ which is equal to $\ell_x(\alpha)$.

2.2. Free factor and free splitting graphs

There are several analogues of the curve graph in the setting of $\text{Out}(\mathbb{F}_n)$. The two important ones are the *free factor graph* $\mathcal{F}(\mathbb{F}_n)$ and the *free splitting graph* $\mathcal{S}(\mathbb{F}_n)$.

A *free factor* \mathcal{A} of \mathbb{F}_n is a subgroup such that there exists another subgroup \mathcal{B} where

$$\mathbb{F}_n = \mathcal{A} * \mathcal{B}.$$

A free factor \mathcal{A} is *proper* in \mathbb{F}_n if the rank of \mathcal{A} is strictly less than n . Then $\mathcal{F}(\mathbb{F}_n)$ is a graph whose vertices are conjugacy classes of proper free factors of \mathbb{F}_n and edges are associated to pairs of proper free factors where one is contained in the other. Similarly for each free factor

\mathcal{A} , one defines the free factor graph $\mathcal{F}(\mathcal{A})$ whose vertices are conjugacy classes of proper free factors of \mathcal{A} .

A *free splitting* over \mathbb{F}_n is a minimal, simplicial (but possibly not free) action of the group \mathbb{F}_n on a simplicial tree T with trivial edge stabilizer. Then $\mathcal{S}(\mathbb{F}_n)$ is a graph whose vertices are free splittings of \mathbb{F}_n up to an equivariant isometry and two splittings are connected by an edge if one can be obtained from the other by a collapse map (see [19] for more details). As above, for any free factor \mathcal{A} , $\mathcal{S}(\mathcal{A})$ denotes the free splitting graph of \mathcal{A} .

There is a projection map

$$\pi: \mathcal{S}(\mathbb{F}_n) \rightarrow \mathcal{F}(\mathbb{F}_n)$$

defined as follows. Let α be a *primitive* loop, that is, $\alpha = [w]$ and $\langle w \rangle$ is a free factor. Then, for any free splitting T , the translation length of α in T , $\ell_T(\alpha)$, can be defined as before but may be zero. We define $\pi(T) = \alpha$ where α is the primitive loop with shortest translation length. Similarly, there is a projection map $\mathcal{S}(\mathcal{A}) \rightarrow \mathcal{F}(\mathcal{A})$ which we also denote by π . These maps are coarsely well defined (see [5]) and further we have

THEOREM 2.1 [20]. *The projection under π of a quasi-geodesic in $\mathcal{S}(\mathcal{A})$ is a reparameterized quasi-geodesic in $\mathcal{F}(\mathcal{A})$.*

REMARK 2.2. The above theorem implies that, given a path p in $\text{Out}(\mathbb{F}_n)$, if its projection to the free factor graph is not a quasi-geodesic, then its projection to the free splitting graph was also not a quasi-geodesic.

Let x be a labelled graph and T be the universal cover. For a proper free factor \mathcal{A} let $T|\mathcal{A}$ be the minimal \mathcal{A} -invariant subtree of T . Note that $T|\mathcal{A} \in \mathcal{S}(\mathcal{A})$. Also recall that the Outer Space (see [11] and the discussion in Section 5) denoted CV_n , is the set of volume 1 metric graphs whose fundamental group is identified with \mathbb{F}_n . Letting T_0 be the universal cover R_0 , the labelled rose fixed in the introduction, we define a shadow map as follows:

$$\Theta_{\mathcal{A}}^{\mathcal{S}}: \text{Out}(\mathbb{F}_n) \rightarrow \mathcal{S}(\mathcal{A}) \quad \phi \rightarrow \phi(T_0)|\mathcal{A}.$$

That is, we change the action on T_0 according to ϕ and take the minimal \mathcal{A} -invariant subtree. Composing with $\pi: \mathcal{S}(\mathcal{A}) \rightarrow \mathcal{F}(\mathcal{A})$ we can define a shadow map to $\mathcal{F}(\mathcal{A})$:

$$\Theta_{\mathcal{A}}^{\mathcal{F}}: \text{Out}(\mathbb{F}_n) \rightarrow \mathcal{F}(\mathcal{A}) \quad \phi \rightarrow \pi(\Theta_{\mathcal{A}}^{\mathcal{S}}(\phi)).$$

The map $\Theta_{\mathcal{A}}^{\mathcal{F}}$ is used more often and hence we shorten the notation to $\Theta_{\mathcal{A}}$. In fact, the shadow to $\mathcal{F}(\mathcal{A})$ makes sense for all marked graphs. That is, we can define a map:

$$\Theta_{\mathcal{A}}^*: \text{CV}_n \rightarrow \mathcal{F}(\mathcal{A}) \quad \Theta_{\mathcal{A}}^*(x) := \pi(T|\mathcal{A}),$$

where T is the universal cover of x . Note that since the image of $\Theta_{\mathcal{A}}^*$ does not depend on the metric of the graph in question, it makes sense to write $\Theta_{\mathcal{A}}^*(x)$ even when x is a marked graph with no metric.

We recall the lemma in [5] on the upper bound for the distance in the free factor graph. It states that if a loop $\alpha \in \mathcal{A}$ is short in x then the shadow of x in $\mathcal{F}(\mathcal{A})$ is near α .

LEMMA 2.3 [5, Lemma 3.3]. *Let \mathcal{A} be a proper free factor and α be a primitive element in \mathcal{A} . Let x be a marked graph so that $\ell_x(\alpha) \leq L$. Then*

$$d_{\mathcal{F}(\mathcal{A})}(\Theta_{\mathcal{A}}^*(x), \alpha) \leq 6L + 13.$$

2.3. *Intersection core*

To define the relative twisting, we first need to introduce the Guirardel Core associated to a pair of \mathbb{F}_n -trees. We give a characterization of the 2-skeleton of the Guirardel Core that is different from (but equivalent to) the one given in [18]. Given an \mathbb{F}_n -tree T , let \mathfrak{o} be a fixed vertex of T that we call the base point. We refer to the vertex $w(\mathfrak{o})$ of T , $w \in \mathbb{F}_n$, simply by w_T and we refer to an edge by a pair of words $(w, ws)_T$ where s is the *label* of the oriented edge. We say the edge $(w, ws)_T$ is an *s-edge*. We say an *s-edge* $(w, ws)_T$ is *preceded* by a *t-edge* $(u, ut)_T$ if $w = ut$.

There is a one-to-one correspondence between the set of infinite geodesic rays originating from \mathfrak{o} and the set of infinite freely reduced words in \mathbb{F}_n . Hence, the Gromov boundary of the group $\partial\mathbb{F}_n$, that is, the equivalence class of quasi-geodesics rays in a Cayley graph \mathbb{F}_n , can be identified with the set of all geodesic rays starting from \mathfrak{o} which in turn can be identified with the Gromov boundary ∂T of T . An (oriented) edge $e = (w, ws)_T$ in T defines a decomposition of $\partial\mathbb{F}_n$ into two sets in the following way: a vertex v_T in T is in *front* of e if the geodesic connecting v_T and $(ws)_T$ does not contain the edge e ; likewise a vertex v_T is *behind* e if the geodesic connecting v_T and w_T does not contain e . Let $\partial^+(e)$ be the set of all the geodesic rays originating from \mathfrak{o} that eventually lie in front of e and $\partial^-(e)$ be set of all geodesic rays originating from \mathfrak{o} that eventually lie behind e .

Note that the sets $\partial^\pm(e)$ are independent of the choice of \mathfrak{o} . Also,

$$\partial^+(e) \cup \partial^-(e) = \partial\mathbb{F}_n \quad \text{and} \quad \partial^+(e) \cap \partial^-(e) = \emptyset.$$

That is, e induces a partition of $\partial\mathbb{F}_n$.

DEFINITION 2.4. Consider an edge e_1 in T_1 and an edge e_2 in T_2 . We say e_1 and e_2 are *boundary equivalent* if they induce the same partition of $\partial\mathbb{F}_n$. We say $e_1 \times e_2$ is an *intersection square* if all of the following four intersections, as subsets of $\partial\mathbb{F}_n$, are nonempty:

$$\begin{aligned} \partial^+(e_1) \cap \partial^+(e_2) &\neq \emptyset & \partial^+(e_1) \cap \partial^-(e_2) &\neq \emptyset \\ \partial^-(e_1) \cap \partial^+(e_2) &\neq \emptyset & \partial^-(e_1) \cap \partial^-(e_2) &\neq \emptyset. \end{aligned}$$

It follows from definition that a pair of boundary equivalent edges never form an intersection square.

EXAMPLE 2.5. Consider the pair of trees T_1 and T_2 that are covers of roses R_1 and R_2 with edge labels $\langle a, a^2b \rangle$ and $\langle a, b \rangle$, respectively. Let \mathfrak{o}_1 be the base point in T_1 and \mathfrak{o}_2 be the base points in T_2 (Figure 2).



FIGURE 2 (colour online). The edges e_1 in T_1 and e_2 in T_2 form an intersection square.

We check that the edge $e_1 = (\text{id}, a)_{T_1}$ and $e_2 = (a, a^2)_{T_2}$ form an intersection square by producing the beginning of the infinite reduced words that lie in each of the intersections:

$$\begin{aligned} \partial^+(e_1) \cap \partial^+(e_2) &: a^2b^{-1} \dots & \partial^+(e_1) \cap \partial^-(e_2) &: a(a^2b)^{-1}aa \dots = ab^{-1} \dots \\ \partial^-(e_1) \cap \partial^+(e_2) &: a^2b \dots & \partial^-(e_1) \cap \partial^-(e_2) &: b^{-1} \dots \end{aligned}$$

Intuitively, this is because the edge $(\text{id}, a^2b)_{T_1}$ in T_1 that is behind e_1 maps to an edge path that passes e_2 , but e_1 maps to an edge that is behind e_2 . Therefore, e_1 and e_2 are tangled to each other.

Let $\overline{\text{Core}(T_1, T_2)}$ be the sub-complex of $T_1 \times T_2$ that is the union of all intersection squares:

$$\overline{\text{Core}(T_1, T_2)} = \{e_1 \times e_2 \mid e_i \in T_i, \quad e_1 \times e_2 \text{ is an intersection square}\}.$$

We define the *intersection core*, or simply the *core* of T_1, T_2 , to be

$$\text{Core}(T_1, T_2) = \overline{\text{Core}(T_1, T_2)} / \mathbb{F}_n.$$

It follows from [9, Lemma 3.4] that the above definition is the same as the definition given by Guirardel.

It is clear from the definition that the core is symmetric: $\text{Core}(T_1, T_2)$ is isomorphic to $\text{Core}(T_2, T_1)$. For an edge $e_2 \in T_2$, the e_2 -slice of the intersection core is the subtree in T_1 that is a collection of edges that form intersection squares with e_2 :

$$C_{e_2}(T_1) = \left\{ e \in T_1 \mid e \times e_2 \text{ is a square in } \overline{\text{Core}(T_1, T_2)} \right\}.$$

Guirardel showed [18] that the e_2 -slice is always convex and finite. That is, if two edges $e, e' \in T/\mathbb{F}_n$ are in a given slice, then all the edges on the geodesic path in T_1/\mathbb{F}_n connecting e and e' are also in the slice.

2.4. Stallings folding path

We describe Stallings folding path here as needed in this paper; for full generality, see [30]. Let x be a labelled directed graph. Then the edges of x can be subdivided into *edgelets* where each edgelet is labelled with an element in the fixed basis $\{s_1, \dots, s_n\}$ of \mathbb{F}_n and the concatenation of edgelets into an edge yields the original labelling of the edge. A *marking-directed fold* from x to x' is a map from x to x' that identifies two edgelets e_1, e_2 for which both of the following are satisfied:

- (i) e_1 and e_2 share the same origin vertex;
- (ii) e_1 and e_2 shares the same label

The resulting quotient map from x to x' is a homotopy equivalence respecting the markings. We recall Stallings' folding theorem per the context of this paper [30].

THEOREM 2.6 Stallings' folding theorem [30]. *For any labelled graph x , there exists a finite sequence of marking-directed folds $x = x_m \rightarrow \dots \rightarrow x_1 \rightarrow x_0 = R_0$ connecting x to R_0 .*

For the labelled rose R in Theorem A, there is a unique Stallings folding path connecting R to R_0 because at every step, there are only two edgelets with the same label and the same original vertex.

2.5. Folding for roses

Here, describe a folding map and its association to partitions of $\partial\mathbb{F}_n$ which is similar to marking-directed fold but one does not need to match the labels of the edges identified. For simplicity, we restrict our attention to cases where the labelled graph is a rose which is sufficient for all our examples.

DEFINITION 2.7. Consider a labelled rose R and choose two edges of R labelled s and t . Let R' be a labelled rose whose edge labels are the same as R except the edge label t has changed to $\bar{s}t$. We say R' is obtained from R by a *fold*, and write

$$R' = \text{fold}(R, t, s) \quad \text{or simply} \quad R' = \text{fold}(R).$$

We also write $T' = \text{fold}(T, t, s)$ for the equivariant map in the universal covers.

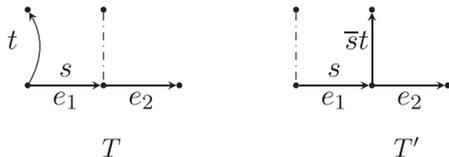


FIGURE 3. An edge labeled t is mapped to two edges labeled s and $\bar{s}t$.

REMARK 2.8. In this paper, a fold can occur between any two edges. In contrast, marking-directed folds [30] follow the labelling of the graphs.

Given \mathbb{F}_n -trees T, T' , and a fixed base points $\mathfrak{o} \in T$ and $\mathfrak{o}' \in T'$, there is a natural morphism $f: T \rightarrow T'$ constructed as follows. Send \mathfrak{o} to \mathfrak{o}' and, for $w \in \mathbb{F}_n$, send the vertex w_T in T to the vertex $w_{T'}$ in T' . Also, send an edge $(w, ws)_T$ to the unique embedded edge path connecting $w_{T'}$ to $(ws)_{T'}$ in T' . The morphism f also induces an \mathbb{F}_n -equivariant homeomorphism

$$f_\infty: \partial T \rightarrow \partial T'.$$

Let T and T' be the universal covers of R and R' respectively. A fold from R to R' induces a morphism from T to T' where an edge of the form $(w, wt)_T$ is mapped to the edge path

$$[(w, ws)_{T'}, (ws, ws(\bar{s}t))_{T'}]$$

and every other edge is mapped to a single edge. Similarly, the edge $(w, ws)_{T'} \in T'$ has two pre-images, $(w, ws)_T$ and a half of $(w, wt)_T$. All other edges have exactly one pre-images. We now describe how partitions given by edges in T differ from that of edges in T' .

PROPOSITION 2.9. Let $R' = \text{fold}(R, t, s)$ and $f: T \rightarrow T'$ be the above morphism. Then

- (i) if an edge e in T is not an s -edge or a t -edge then $f(e) \in T'$ is boundary equivalent to e ;
- (ii) if $e = (w, wt)_T$ is a t -edge, then $e' = (ws, ws(\bar{s}t))_{T'}$, which is contained in $f(e)$, is boundary equivalent to e ;
- (iii) if e_2 is an s -edge and e_1 is the t -edge starting at the same vertex as e_2 then $e' = f(e_2)$ partitions $\partial\mathbb{F}_n$ in the following way:

$$\partial^-(e') = \partial^-(e_2) \setminus \partial^+(e_1) \quad \text{and} \quad \partial^+(e') = \partial^+(e_2) \cup \partial^+(e_1);$$

- (iv) for adjacent edges $e_1 = (u, us)_T, e_2 = (us, w)_T$ and $e' \subset f(e_2)$, we have

$$\partial^-(e_1) \subset \partial^-(e').$$

Proof. Let \mathfrak{o} and \mathfrak{o}' be the base points in T and T' ; $F(\mathfrak{o}) = \mathfrak{o}'$. Note that changing a base point \mathfrak{o} to a point w in T moves both the partition given by e in T and given by e' in T' by the action of w . Hence, it is sufficient to prove the statement for any desired base point.

Consider an edge $e' \in T'$ that has only one pre-image under the morphism $f: T \rightarrow T'$. That is, e is the only edge where $e' \subseteq f(e)$. Then, e and e' are boundary equivalent. To see this, assume $\mathfrak{o} = w_T$. A ray r starting from \mathfrak{o} crosses e if and only if $f(r)$ crosses e' . Hence, a ray is

eventually in front of e if and only if it is eventually in front of e' which, by definition, means e and e' are boundary equivalent. This finished the proof of the first two parts.

Let $e_1 = (u, ut)_T$ and $e_2 = (u, us)_T$. We choose $\mathfrak{o} = u$. Then e_1 and e_2 are the only two edges that are mapped over $f(e_1)$. But a ray r in T starting from \mathfrak{o} crosses at most one of e_1 or e_2 . In fact, $f(r)$ crosses either edge if and only if it is eventually in front of $f(e_1)$. That is,

$$\partial^+ f(e) = \partial^+(e) \cup \partial^+(\hat{e}).$$

The other equality in part (iii) holds because $\partial^- f(e_1)$ is the complement of $\partial^+ f(e_1)$.

In part (iv), we have $e_1 = (u, us)_T$. Since $R' = \text{fold}(R, t, s)$, the only ray r in $\partial^-(e_1)$ where $f(r)$ crosses $f(e_1)$ is a ray starting with $(u, ut)_T$. If such is the case, then $f(r)$ starts with $[(u, us)_{T'}, (us, us(\bar{s}t))_{T'}]$. But e' is different from $(us, us(\bar{s}t))_{T'}$ (which has only one pre-image). Hence, $f(r)$ does not cross e' and lands in $\partial^+(e')$. Therefore,

$$\partial^-(e_1) \subset \partial^-(e'). \quad \square$$

3. Twisting estimate

Let $|G|$ denote the number of edges in the given graph G . We now define *relative twisting number*, which is an analogue of the Masur–Minsky twisting number [26]. Recall that, for a given pair of trees T and T_0 , a slice $C_{e_0}(T)$ over an edge $e_0 \in T_0$ is the subtree of T consisting of edges that form intersection squares with e_0 .

DEFINITION 3.1. Given a loop α and two \mathbb{F}_n -trees T, T_0 , the *relative twisting number* of \mathbb{F}_n -trees T, T_0 around α is

$$\text{twist}_\alpha(T, T_0) = \max_{e_0 \in T_0, w \in \alpha} \frac{|\text{axis}_T(w) \cap C_{e_0}(T)|}{\ell_T(\alpha)}.$$

If T and T_0 are universal covers of R and R_0 , we define $\text{twist}_\alpha(R, R_0) = \text{twist}_\alpha(T, T_0)$.

EXAMPLE 3.2. We illustrate the intuition behind this definition by an example. Consider the rose R with labels $\langle a, b, c(ab)^4 \rangle$ and let T be the universal cover of R with base point \mathfrak{o} and let T_0 be the universal cover of R_0 with labels $\langle a, b, c \rangle$ and base point \mathfrak{o}_0 as usual. Let $\alpha = [ab]$ and $w = ab$. Then, $\ell_T(\alpha) = 2$. Consider the edges $e_a = (\text{id}, a)_{T_0}$. Then C_{e_a} is the tree depicted next. The solid edges form a slice but the dashed edges are not in the core. One can see this by direct computation or from [9, Lemma 3.7] which states that C_{e_a} is the interior of the convex hull of the pre-image of e_a . Intuitively, this is caused by the fact that the edge $(\bar{c}, (ab)^4)_T$ in T (which is the dotted edge on the right) is in front of all the solid edges in T but maps to an edge path that crosses e_a which is behind the image of the solid edges in T_0 . The edge $e = (\text{id}, a)_T$ (the dotted edge in the left) does not form an intersection square with e_a because $\partial^- e \cap \partial^+ e_a = \emptyset$ (Figure 4).

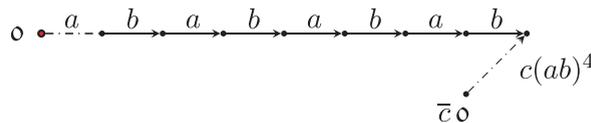


FIGURE 4 (colour online). The slice C_{e_a} over $e_a \in T_0$ is a path of length 7.

The whole slice lies on $\text{axis}_T(ab)$, therefore

$$\frac{|\text{axis}_T(ab) \cap C_{e_a}(T)|}{\ell_T(\alpha)} = \frac{7}{2} = 3.5.$$

Meanwhile, if we consider $e_b = (\text{id}, b)_T$, the slice over e_b is the set of solid edges in Figure 5.

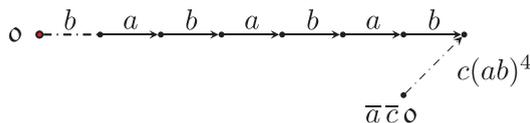


FIGURE 5 (colour online). The slice C_{e_b} over $e_b \in T_0$ is a path of length 6.

Since this slice also lies on $\text{axis}_T(ab)$, we have

$$\frac{|\text{axis}_T(ab) \cap C_{e_b}(T)|}{\ell_T(\alpha)} = \frac{6}{2} = 3.$$

This means, if one starts from b , one observes less twisting. Lastly, the edge $e_c = (\text{id}, c)_T$ has only one pre-image in T . This implies that $C_{e_c} = \emptyset$ (again, see [9, Lemma 3.7]). Thus the estimate for the twist depends heavily on the choice of e_0 . In our example,

$$\text{twist}_\alpha(T, T_0) = \max\{3.5, 3, 0\} = 3.5 \quad \text{and} \quad \text{tw}_\alpha(T, T_0) = [3.5] = 3.$$

REMARK 3.3. The twisting number defined above is a rational number. The integer part of $\text{twist}_\alpha(T, T_0)$, which we denote by $\text{tw}_\alpha(T, T_0)$, is equal to the Clay–Pettet definition of relative twisting number which considers hyperplanes τ, τ_0 in the core that are dual to specific edges and counts how many α -translates of τ intersects τ_0 (see [8] for details). Similar to the above example, one can see that in the example of Theorem A,

$$\text{tw}_b(R_0, \phi(R_0)) = s - 1 \quad \text{and} \quad \text{tw}_{ab^s}(R_0, \phi(R_0)) = t - 1.$$

We now show that the relative twisting number changes slowly along loops with large lengths. For a real number $r > 0$, let $[r]$ be the integer part of r and $\{r\}$ be the fractional part of r .

THEOREM 3.4. Let R' be obtained from R by a single fold:

$$R' = \text{fold}(R, t, s).$$

Then, for any loop α ,

$$\text{tw}_\alpha(R', R_0) \geq \text{tw}_\alpha(R, R_0) - 1, \tag{1}$$

where R_0 is as defined before. Furthermore,

$$\text{twist}_\alpha(R', R_0) \geq \left[\text{twist}_\alpha(R, R_0) - \frac{2}{\ell_R(\alpha)} \right] + \frac{\left\{ \text{twist}_\alpha(R, R_0) - \frac{2}{\ell_R(\alpha)} \right\}}{4}. \tag{2}$$

We need to prepare for the proof by establishing a few lemmas. Note that R can also be obtained from R' by a single fold:

$$R' = \text{fold}(R, t, s), \quad R = \text{fold}(R', \bar{s}t, \bar{s}).$$

For the rest of this section, we assume T and T' are universal covers of R and R' respectively and that

$$f: T \rightarrow T' \quad \text{and} \quad g: T' \rightarrow T$$

are the morphism associated to these folds. For an embedded edge path $P = [e_1 \cdots e_k]$ in T , let $f(P)$ denote the image of P under the morphism f and let $f(P)_w$ denote the embedded edge path that is the intersection of $f(P)$ and $\text{axis}_{T'}(w)$. We call e_1 and e_k the end edges of P .

LEMMA 3.5. *For any loop α , we have*

$$\ell_T(\alpha) \geq \frac{\ell_{T'}(\alpha)}{2}.$$

Also, if $P = [e_1 \cdots e_k]$ is an edge path on the axis $_T(w)$, for some $w \in \mathbb{F}_n$, so that both $f(e_1)$ and $f(e_k)$ contain an edge on axis $_{T'}(w)$, then

$$|f(P)_w| \geq \frac{|P|}{2}.$$

Proof. Recall that f maps an edge $(u, ut)_T$ to the edge path

$$[(u, us)_{T'}, (us, us(\bar{st}))_{T'}]$$

and maps every other edge to one edge. Therefore, for any embedded edge path P in T ,

$$|f(P)| \leq 2|P|. \quad (3)$$

If P is an edge path that realizes $\ell_T(\alpha)$, then $|f(P)| \geq \ell_{T'}(\alpha)$, thus

$$\ell_T(\alpha) = |P| \geq \frac{|f(P)|}{2} \geq \frac{\ell_{T'}(\alpha)}{2}.$$

Now assume $P = [e_1 \cdots e_k]$ is an edge path on the axis $_T(w)$ and both $f(e_1)$ and $f(e_k)$ contain an edge on axis $_{T'}(w)$. Applying, equation (3) to the morphism g and the edge path $f(P)_w$, we have

$$|g(f(P)_w)| \leq 2|f(P)_w|. \quad (4)$$

We need to show $P \subseteq g(f(P)_w)$. In fact, it suffices to show that any end vertex of P is contained in $g(f(P)_w)$.

Let u_T be the first vertex in e_1 (the argument for the last vertex of e_k is similar). If $u_{T'} \in \text{axis}_{T'}(w)$, then $u_{T'} \in f(P)_w$, which means $u_T \in g(f(P)_w)$. Otherwise, $u_{T'} \notin \text{axis}_{T'}(w)$. By assumption, $f(e_1)$ contains an edge on axis $_{T'}(w)$. It follows that e_1 is mapped to two edges $[e', e'']$ under f , which means e_1 is necessarily a t -edge (or a t^{-1} -edge, in which the remainder of the proof changes accordingly). Furthermore, the edge e' has two pre-images. This is because, $u_{T'} \notin \text{axis}_{T'}(w)$ and the edge preceding e_1 along the axis $_T(w)$ must be mapped over e' as well. But the edges in T' with label \bar{st} have only one pre-image. Hence,

$$e_1 = (u, ut)_T, \quad e' = (u, us)_{T'}, \quad \text{and} \quad e'' = (us, ut)_{T'}.$$

We then have

$$(us)_{T'}, (ut)_{T'} \in \text{axis}_{T'}(w) \implies (us)_T, (ut)_T \in g(f(P)_w).$$

But T is a tree and the vertex u_T necessarily lies on the path connecting $(us)_T$ and $(ut)_T$. Thus $(us)_T, (ut)_T \in g(f(P)_w)$ implies $u_T \in g(f(P)_w)$.

We have shown that the end vertices of P are both in $g(f(P)_w)$, which implies $P \subseteq g(f(P)_w)$. The lemma follows from equation (4). \square

Let $\alpha \in \mathbb{F}_n$ be a loop. Let $e_0 \in T_0$ and $w \in \alpha$ be so that the edge path

$$P := [e_1 e_2 \dots e_k] = \text{axis}_T(w) \cap C_{e_0}(T)$$

is the one realizing the maximum in the definition of $\text{twist}_\alpha(T, T_0)$, although we remark that maximality is not needed for the following result.

LEMMA 3.6. *Let P be the path above and assume $|P| \geq 2$. Then $f([e_1 e_2])$ contains an edge in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$. The same holds for $f([e_{k-1} e_k])$. Furthermore, if $\ell_T(\alpha) = 1$, then either $f(e_1)$ or $f(e_k)$ contains an edge in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$.*

Proof. The edge path $f(P)$ contains $f(P)_w$. Furthermore, an edge in $f(P)$ is on $f(P)_w$ if and only if it has a unique pre-image in $\text{axis}_T(w)$.

If, for $i = 1$ or 2 , $e_i = (u, v)$ and $\bar{u}v$ is not s, \bar{s}, t or \bar{t} , then $f(e_i)$ is a single edge, has one pre-image and it is boundary equivalent to e_i . Hence, it lies on $C_{e_0}(T') \cap \text{axis}_{T'}(w)$.

If $e_i = (u, ut)$, by the definition of $\text{fold}(T, t, s)$, e_i is mapped to the edge path

$$[(u, us)(us, us(\bar{s}t))].$$

Since the edge $(us, us(\bar{s}t))$ has only one pre-image, it lies on $f(P)_w$ and by Proposition 2.9, it is boundary equivalent to e_i and thus $(us, us(\bar{s}t)) \in C_{e_0}(T')$. Similar argument works when $e_i = (ut, u)$. Also, this shows that if $\alpha = t$ or \bar{t} , then either $f(e_1)$ or $f(e_k)$ contains an edge in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$.

There are two remaining cases. Assume $e_1 = (u, us)$ and $e_2 = (us, us^2)$. Then the t -edge starting at us , (us, ust) , does not lie on $\text{axis}_T(w)$. Hence, $f(e_2)$ has one pre-image in $\text{axis}_T(w)$ and thus is on $\text{axis}_{T'}(w)$. Also, by part (iii) of Proposition 2.9,

$$\partial^+(e_2) \subset \partial^+(f(e_2))$$

and by part (iv) of Proposition 2.9,

$$\partial^-(e_1) \subset \partial^-(f(e_2)).$$

Since $e_1, e_2 \in C_{e_0}(T)$, each of $\partial^+(e_2)$ and $\partial^-(e_1)$ intersects each of $\partial^+(e_0)$ and $\partial^-(e_0)$. Therefore, each of $\partial^+(f(e_2))$ and $\partial^-(f(e_2))$ intersects each of $\partial^+(e_0)$ and $\partial^-(e_0)$. That is, $f(e_2) \in C_{e_0}(T')$.

The remaining case when $e_1 = (us^2, us)$ and $e_2 = (us, u)$ is identical, except in this case, $f(e_1)$ is in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$. These two cases also show that if $\alpha = s$ or \bar{s} , then both $f(e_1)$ and $f(e_k)$ contains an edge in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$. \square

Proof of Theorem 3.4. Recall $\alpha \in \mathbb{F}_n$ is a loop and the edge path

$$P := [e_1 e_2 \dots e_k] = \text{axis}_T(w) \cap C_{e_0}(T)$$

is the edge path that realizes the maximum in the definition of $\text{twist}_\alpha(T, T_0)$. By Lemma 3.6, either $f(e_1)$ or $f(e_2)$ contains an edge that is in $C_{e_0}(T') \cap \text{axis}_{T'}(w)$. Call the associated edge in T (either e_1 or e_2), e_{first} . Likewise, one of $f(e_{k-1})$ or $f(e_k)$ has this property. Call the associated edge in T (either e_{k-1} or e_k) e_{last} . The combinatorial length of the edge path $[e_{\text{first}} \dots e_{\text{last}}]$ is at least

$$k - 2 = \text{twist}_\alpha(T, T_0) \ell_T(\alpha) - 2.$$

Hence, if we define

$$p := \left\lceil \text{twist}_\alpha(T, T_0) - \frac{2}{\ell_R(\alpha)} \right\rceil, \tag{5}$$

then $(k - 2) \geq p \cdot \ell_R(\alpha) = p \cdot \ell_T(\alpha)$. Therefore, $[e_{\text{first}} \dots e_{\text{last}}]$ contains at least p copies of the fundamental domain of the action of w on $\text{axis}_T(w)$ and, in particular, $w^p(e_{\text{first}})$ (the translation of e_{first} by the action of w^p) lies on $[e_{\text{first}} \dots e_{\text{last}}]$. Let P_{remain} be the edge path $[w^p(e_{\text{first}}) \dots e_{\text{last}}]$. Since e_{first} contains an edge in $\text{axis}_{T'}(w)$, so does $w^p(e_{\text{first}})$. That is, $f(P)_w$ consists of p fundamental domains of the action of w on $\text{axis}_{T'}(w)$ and the segment $f(P_{\text{remain}})_w$. Thus

$$\text{twist}_\alpha(T', T_0) \geq \frac{p \ell_{T'}(\alpha) + |f(P_{\text{remain}})_w|_{T'}}{\ell_{T'}(\alpha)} = p + \frac{|f(P_{\text{remain}})_w|_{T'}}{\ell_{T'}(\alpha)}. \tag{6}$$

Also, by Lemma 3.5

$$\ell_{T'}(\alpha) \leq 2 \ell_T(\alpha) \quad \text{and} \quad |f(P_{\text{remain}})_w|_{T'} \geq \frac{|P_{\text{remain}}|_T}{2}.$$

Also, since

$$|P_{\text{remain}}|_T \geq |P|_T - p \ell_T(\alpha) \geq (k-2) - p \ell_T(\alpha),$$

we have

$$\frac{|P_{\text{remain}}|_T}{\ell_T(\alpha)} \geq \frac{k}{\ell_T(\alpha)} - \frac{2}{\ell_T(\alpha)} - p.$$

Let $L := \ell_T(\alpha)$, we then have

$$\begin{aligned} \frac{|P_{\text{remain}}|_T}{\ell_T(\alpha)} &\geq \frac{k}{\ell_T(\alpha)} - \frac{2}{\ell_T(\alpha)} - p \\ &\geq \text{twist}_\alpha(T, T_0) - \frac{2}{L} - p = \left\{ \text{twist}_\alpha(T, T_0) - \frac{2}{L} \right\}, \end{aligned}$$

where the last equality follows from the definition of p . Hence

$$\frac{|f(P_{\text{remain}})_w|_{T'}}{\ell_{T'}(\alpha)} \geq \frac{|P_{\text{remain}}|_T/2}{2\ell_T(\alpha)} \geq \frac{\left\{ \text{twist}_\alpha(T, T_0) - \frac{2}{L} \right\}}{4}. \quad (7)$$

Equation (2) follows from equations (5)–(7).

For $L \geq 2$, equation (2) implies equation (1). If $L = 1$, from the second assertion of Lemma 3.6, we have that either $e_1 = e_{\text{first}}$ or $e_k = e_{\text{last}}$. Hence, $f(P)_w$ contains at least $(k-1)$ fundamental domains of action of w and $\text{tw}_\alpha(T', T_0) \geq k-1$. But $\text{tw}_\alpha(T, T_0) = k$. Therefore, equation (1) still holds. \square

THEOREM 3.7. *For any folding sequence T_m, \dots, T_0 and any loop α , we have*

$$m \geq \text{tw}_\alpha(T_m, T_0).$$

Further, if $\ell_{T_i}(\alpha) \geq L > 50$ for every i , then

$$m > \text{tw}_\alpha(T_m, T_0) \left(\log_5 \frac{L}{50} \right).$$

Proof. The first assertion of the theorem follows directly from the first assertion of Theorem 3.4. We prove the second assertion.

For a real number $\frac{10}{L} < r$, we have

$$r - \frac{2}{L} > \frac{4r}{5} \quad \text{and} \quad \frac{r - \frac{2}{L}}{4} > \frac{r}{5}. \quad (8)$$

For any $0 \leq N < \text{tw}_\alpha(T_m, T_0)$, consider the first time $i = i(N)$ when $\text{tw}_\alpha(T_i, T_0) = N$. Since the index decreases as we fold towards T_0 , i is the maximal of all the times $\text{tw}_\alpha(T_j, T_0) = N$. Now apply Theorem 3.4 with $T = T_{i+1}$ and $T' = T_i$. Since i is the first time $\text{tw}_\alpha(T_i, T_0) = N$, we have

$$[\text{twist}(T_i, T_0)] = N \quad \text{and} \quad [\text{twist}(T_{i+1}, T_0)] = N + 1.$$

That is to say, $[\text{twist}(T_{i+1}, T_0) - \frac{2}{L}] \geq N$ as $\frac{2}{L} < 1$. By equation (2),

$$\text{twist}(T_i, T_0) \geq \left[\text{twist}(T_{i+1}, T_0) - \frac{2}{L} \right] + \frac{\left\{ \text{twist}(T_{i+1}, T_0) - \frac{2}{L} \right\}}{4}$$

therefore

$$N + \{\text{twist}(T_i, T_0)\} \geq N + \frac{\left\{ \text{twist}(T_{i+1}, T_0) - \frac{2}{L} \right\}}{4}$$

and hence

$$\{\text{twist}(T_i, T_0)\} \geq \frac{\{\text{twist}(T_{i+1}, T_0) - \frac{2}{L}\}}{4}.$$

Again since this is the first time that $\text{tw}_\alpha(T_i, T_0) = N$, we have

$$\text{twist}(T_{i+1}, T_0) \geq N \geq 0.$$

That is to say $\text{twist}(T_{i+1}, T_0) \geq 1$. From equation (2), we also have

$$\text{twist}_\alpha(T_i, T_0) \geq \left[\text{twist}_\alpha(T_{i+1}, T_0) - \frac{2}{L} \right].$$

Combined with the fact that $\text{tw}_\alpha(T_i, T_0) < \text{tw}_\alpha(T_{i+1}, T_0)$, we have that

$$\{\text{twist}_\alpha(T_{i+1}, T_0)\} < \frac{2}{L}. \quad (9)$$

Therefore,

$$\begin{aligned} \{\text{twist}(T_i, T_0)\} &\geq \frac{\{\text{twist}(T_{i+1}, T_0) - \frac{2}{L}\}}{4} \\ &= \frac{\{\text{twist}(T_{i+1}, T_0) + 1 - \frac{2}{L}\}}{4} \\ &= \frac{\{\{\text{twist}(T_{i+1}, T_0)\} + 1 - \frac{2}{L}\}}{4} \\ &\geq \frac{\{1 - \frac{2}{L}\}}{4} = \frac{1 - \frac{2}{L}}{4} \geq \frac{1}{5}. \end{aligned}$$

Let $(i - k - 2)$ be first time when $\text{tw}_\alpha(T_{i-k-1}, T_0) = N - 1$. Then the step immediately before that we again have

$$\{\text{twist}(T_{i-k-1}, T_0)\} \leq \frac{2}{L}.$$

Which means for the step before that we have

$$\frac{\{\text{twist}(T_{i-k}, T_0) - \frac{2}{L}\}}{4} \leq \{\text{twist}(T_{i-k-1}, T_0)\} \leq \frac{2}{L}.$$

Which means $\{\text{twist}(T_{i-k}, T_0)\} \leq \frac{10}{L}$. That is, in k -steps, the fractional twist has been reduced from above $\frac{1}{5}$ to below $\frac{10}{L}$. Thus, from Proposition 3.4 and equation (8), we have

$$\frac{1}{5^k} < \frac{10}{L},$$

which can be rewritten as

$$k > \log_5 \frac{L}{50}.$$

Since this is true for every $0 \leq N < \text{tw}_\alpha(T_0, T_m)$, we have

$$\text{tw}_\alpha(T_m, T_0) < \frac{m}{\log_5 \frac{L}{50}} \implies m > \text{tw}_\alpha(T_m, T_0) \left(\log_5 \frac{L}{50} \right),$$

which is as desired. \square

3.1. The lower bound is sharp

At first glance, one might think, that the factor $\log_5 \frac{L}{50}$ in Theorem 3.4 could be replaced with a linear function of L . However, we show that the above estimate is sharp up to a uniform

multiplicative error. We now construct, for an arbitrarily large L a folding path $R = R_m, \dots, R_0$ so that

- the length of α at each R_i is at least L ;
- m is comparable with $\log L \operatorname{tw}_\alpha(R, R_0)$.

EXAMPLE 3.8. Let R be a rose of rank 5 with edge labels

$$\langle (bc)^m a, db, \phi^k(d)c, d, e \rangle,$$

where $\phi: \langle d, e \rangle \rightarrow \langle d, e \rangle$ is a fully irreducible automorphism of the free factor $\langle d, e \rangle$ with exponential growth and

$$\lambda(\phi) = \lambda(\phi^{-1}).$$

Let L be the word length of $\phi^{\lfloor k/2 \rfloor}(d)$ (and also the length of $\phi^{-\lfloor k/2 \rfloor}(d)$) in $\langle d, e \rangle$ which can be chosen to be arbitrarily large by choosing k large enough. Let $\alpha = [bc]$. Let R_0 be a rose with labels

$$\langle a, db, \phi^k(d)c, d, e \rangle.$$

The shortest way to express α in R is

$$bc = (\bar{d}) \cdot (db) \cdot \phi^{-k}(d) \cdot (\phi^k(d)c),$$

where the terms in parentheses are labels of edges in R and is a word of length roughly L^2 in $\langle d, e \rangle$. We have

$$\ell_R(\alpha) \asymp L^2 > L.$$

Similar to Example 3.2, we have

$$\operatorname{tw}_\alpha(R, R_0) = m - 1.$$

We now start twisting around the loop $[bc]$, however in a somewhat un-natural way that always keeps the length of α larger than L , using the following steps.

(1) Twist the $((bc)^m a)$ -loop around the first half of α , that is to say, cancel b (which is half of α) from $(bc)^m a$. To do this we first fold the $((bc)^m a)$ -loop around the (d) -loop, and then fold the resulting loop around (db) -loop. Thus this step uses two folds:

$$\overline{(db)} \cdot d \cdot (bc)^m a = c(bc)^{m-1} a.$$

(2) Fold $\langle d, e \rangle$ to $\langle \phi^k(d), \phi^k(e) \rangle$; this takes

$$k \|\phi\|_{\langle d, e \rangle} \asymp \log L$$

many steps. Note that the immersed loop

$$bc = (\bar{d}) \cdot (db) \cdot \phi^{-k}(d) \cdot (\phi^k(d)c),$$

contains both \bar{d} and $\phi^{-k}(d)$. At any point along this folding path, say after ϕ_i has been applied i times, we have

$$\max \left(|\bar{d}|_{\langle \phi^i(d), \phi^i(e) \rangle}, |\phi^{-k}(d)|_{\langle \phi^i(d), \phi^i(e) \rangle} \right) \geq |\phi^{\lfloor k/2 \rfloor}(d)| = L,$$

where $|\cdot|_{\langle \phi^i(d), \phi^i(e) \rangle}$ denotes the word length of an element in the group $\langle d, e \rangle$ in terms of $\phi^i(d)$ and $\phi^i(e)$. Thus the combinatorial lengths of α remains at least L for each i .

(3) Twist around the second half of α , that is, cancel c from $c(bc)^{m-1} a$ in two folds:

$$\overline{(\phi^k(d)c)} \cdot \phi^k(d) \cdot c(bc)^{m-1} a = (bc)^{m-1} a.$$

(4) Fold $\langle \phi^k(d), \phi^k(e) \rangle$ to $\langle d, e \rangle$. Again, the number of steps is comparable to $\log L$ and the length of α remains larger than L .

We now repeat steps 1–4, m -times. Every time the relative twisting around α is reduced by 1. The path has a length of order $m \log L$ as desired.

4. Quasi-geodesics in $\text{Out}(\mathbb{F}_n)$

In this section we use Theorem B to prove Theorem C. Consider a path $p: [0, m] \rightarrow \mathcal{X}$ from an interval $[0, m] \subset \mathbb{R}$ to a metric space \mathcal{X} and, for $i \in [0, m]$, let $p_i = p(i)$. Recall that p is a (K, C) -quasi-geodesic if for any $i, j \in [0, m]$, we have

$$\frac{|i - j| - C}{K} \leq d_{\mathcal{X}}(p_i, p_j) \leq K|i - j| + C.$$

It follows that, any three points in the image of a quasi-geodesic satisfy a coarse-reverse-triangle inequality. Namely, for any $i, j, k \in [0, m]$, $i \leq j \leq k$, we have

$$d_{\mathcal{X}}(p_i, p_j) + d_{\mathcal{X}}(p_j, p_k) \leq K^2 d_{\mathcal{X}}(p_i, p_k) + (K + 2)C. \tag{10}$$

We say p is a reparametrized (K, C) -quasi-geodesic if there is a reparametrization $\rho: [0, m] \rightarrow [0, m]$ so that $p \circ \rho$ is a (K, C) -quasi-geodesic. Since the image of p and $p \circ \rho$ are the same, if p is a reparametrized quasi-geodesic, any three points in its image still satisfy the coarse-reverse-triangle inequality given in equation (10). Often, it is convenient to consider maps from intervals $[0, m]_{\mathbb{Z}}$ in \mathbb{Z} to a metric space \mathcal{X} . Then we say $p: [0, m]_{\mathbb{Z}} \rightarrow \mathcal{X}$ is a reparametrized quasi-geodesic if it is a restriction of a reparametrized quasi-geodesic from $[0, m] \rightarrow \mathcal{X}$.

We can now restate Theorem C explicitly as follows:

THEOREM 4.1 (Quasi-geodesics backtrack in sub-factors). *For given constants $K_1, C_1, K_2, C_2 > 0$ there exists an automorphism, $\phi \in \text{Out}(\mathbb{F}_n)$ such that, for any (K_1, C_1) -quasi-geodesic $p: [0, m] \rightarrow \text{Out}(\mathbb{F}_n)$ with $p(0) = \text{id}$ and $p(m) = \phi$, the shadow $\Theta_{\mathcal{A}} \circ p$ of p in $\mathcal{F}(\mathcal{A})$ is not a (K_2, C_2) -reparameterized quasi-geodesic.*

Proof. Let $\langle a, b, c \rangle$ be a generating set for \mathbb{F}_3 , and let R_0 be a rose with labels $\{a, b, c\}$. Let $\psi \in \text{Out}(\mathbb{F}_3)$ be an automorphism defined as:

$$\begin{cases} a \rightarrow aba \\ b \rightarrow ab \\ c \rightarrow c \end{cases}$$

Given the generating set introduced in the introduction, $\|\psi\| = \|\psi^{-1}\| = 2$ where $\|\cdot\|$ represents the word length. Let $\mathcal{A} = \langle a, b \rangle < \mathcal{F}_3$ be a rank 2 free factor. The automorphism ψ fixes \mathcal{A} . We denote the restriction of ψ to \mathcal{A} by $\psi_{\mathcal{A}}$. Then, $\psi_{\mathcal{A}}$ is an irreducible automorphism and acts loxodromically on the free-factor graph $\mathcal{F}(\mathcal{A})$ of \mathcal{A} . Hence, there exists a constant $c_{\psi} > 0$ so that, for an integer $q > 0$

$$d_{\mathcal{A}}(R_0, \psi^q(R_0)) \geq c_{\psi}q. \tag{11}$$

Let α be the loop represented by the word $\psi^q(a)$ and, for a large integer $t > 0$, let $\phi \in \text{Out}(\mathbb{F}_3)$ be the automorphism that twists the element c around α t -times, namely:

$$\begin{cases} a \rightarrow a \\ b \rightarrow b \\ c \rightarrow c(\psi^q(a))^t \end{cases}$$

We first find an upper-bound for the $\|\phi\|$ by constructing a path connecting identity to ϕ . First apply ψ^q so α is represented by one edge in the rose $\psi^q(R_0)$, then twist c around α t -times, and then apply ψ^{-q} . We have

$$\|\phi\| \leq q(\|\psi\| + \|\psi^{-1}\|) + t = 4q + t.$$

Now consider the (K_1, C_1) -quasi-geodesic $p: [0, m] \rightarrow \text{Out}(\mathbb{F}_3)$ connecting the identity to ϕ . We have

$$m \leq K_1\|\phi\| + C_1 \leq 4K_1q + K_1t + C_1. \tag{12}$$

If the shadow of p to \mathcal{A} is a (K_2, C_2) -reparametrized quasi-geodesic, then the coarse-reverse-triangle inequality (equation (10)) holds. That is, for any index i and $R_i = p(i)(R_0)$, we have

$$d_{\mathcal{A}}(R_0, R_i) + d_{\mathcal{A}}(R_i, \phi(R_0)) \leq (K_2)^2 d_{\mathcal{A}}(R_0, \phi(R_0)) + (K_2 + 2)C_2.$$

But ϕ fixes a and b and thus R_0 and $\phi(R_0)$ have the same projection to \mathcal{A} . Hence,

$$2d_{\mathcal{A}}(R_0, R_i) \leq (K_2 + 2)C_2.$$

Now, using equation (11), we get

$$d_{\mathcal{A}}(R_i, \psi^q(R_0)) \geq d_{\mathcal{A}}(R_0, \psi^q(R_0)) - d_{\mathcal{A}}(R_0, R_i) \geq c_{\psi}q - (K_2 + 2)C_2.$$

By Lemma 2.3, this implies

$$\ell_{R_i}(\alpha) \geq \frac{c_{\psi}q - (K_2 + 2)C_2 - 13}{6} =: L. \tag{13}$$

Now, Theorem 3.7 implies that

$$t = \text{tw}_{\alpha}(R_0, R_m) \leq \frac{m}{\log_5 \frac{L}{50}}$$

and using equation (12) we get

$$t \leq \frac{4K_1q + K_1t + C_1}{\log_5 \frac{L}{50}}.$$

If we choose q large enough so that

$$\log_5 \frac{L}{50} > 2K_1$$

and then choose t large enough so that

$$\frac{t}{2} > \frac{4K_1q + C_1}{\log_5 \frac{L}{50}},$$

we get

$$t \leq \frac{4K_1q + C_1}{\log_5 \frac{L}{50}} + \frac{K_1t}{\log_5 \frac{L}{50}} < \frac{t}{2} + \frac{t}{2} = t$$

which is a contradiction. The contradiction proves that the shadow of p to $\mathcal{F}(\mathcal{A})$ is not a reparametrized (K_2, C_2) -quasi-geodesic. \square

5. Outer Space

Outer Space CV_n is metric space with $\text{Out}(\mathbb{F}_n)$ action defined as an analogue of the Teichmüller space; see [11] for more details. Here, we introduce CV_n briefly and prove Theorem D.

We assume a *graph* is always simple and all vertices have degree 3 or more. A *marked metric graph* (x, f) is a metric graph x together with homotopy equivalence $f: R_0 \rightarrow x$. The space

of all marked metric graphs whose edge lengths sum up to one is called the *Outer Space* [11] and is denoted by CV_n . The group $\text{Out}(\mathbb{F}_n)$ acts on CV_n by pre-composing the marking: for an element $\phi \in \text{Out}(\mathbb{F}_n)$, $\phi(x, f) = (x, f \circ \phi)$.

Note that we can still think of x as a labelled graph. Recall that $\ell_x(\alpha)$ denotes the combinatorial length of α in x . Let $|e|_x$ denote the metric length of an edge e in x and $|\alpha|_x$ the metric length of α in x , which is the metric length of the immersed loop of the representative of α that realizes its combinatorial length in x .

For a fixed $\epsilon > 0$ define the *thick part* of CV_n to be the set of $x \in \text{CV}_n$ such that

$$|\alpha|_x \geq \epsilon \text{ for every nontrivial conjugacy class } \alpha.$$

A map $h: (x, f_x) \rightarrow (y, f_y)$ is a *difference of markings* map if $h \circ f_x \simeq f_y$ (homotopy). We will only consider Lipschitz maps and we denote by L_h the Lipschitz constant of h . In many ways it is natural to consider the (asymmetric) Lipschitz metric on CV_n :

$$d(x, y) := \inf_h \log L_h,$$

where the infimum is taken over all differences of markings maps. We refer the reader to [2, 15] for review for some metric properties of $d(\cdot, \cdot)$. In particular, there always exists a non-unique difference of markings map that realizes the infimum. Since a difference of markings map is homotopic rel vertices to a map that is linear on edges, we also use h to denote the representative that realizes the infimum and is linear on edges and refer to such a map as an *optimal map* from x to y . For this section we always assume h is an optimal difference of markings map. Since h is linear on edges, we define

$$\lambda(e) = \frac{|h(e)|_y}{|e|_x}$$

to be the *stretch factor* of an edge e and

$$\lambda(\alpha) = \frac{|\alpha|_y}{|\alpha|_x}$$

to be the *stretch factor* of a shortest immersed loop that represents α . Define the *tension subgraph*, x_ϕ , or $\text{stretch}(x, y)$, to be the subgraph of x consisting of maximally stretched edges.

Now we restate and prove part (i) of Theorem D:

THEOREM 5.1 (Shadow of a geodesic in CV_n is not a quasi-geodesic in $\text{Out}(\mathbb{F}_n)$). *For given constants K and C , there are points $x, y \in \text{CV}_n$ such that for every geodesic $[x, y]_{\text{CV}_n}$ in CV_n connecting x to y , its image in $\text{Out}(\mathbb{F}_n)$ is not (K, C) -quasi-geodesic.*

Proof. Consider the same example in rank 3 as in Theorem 4.1 where ϕ is defined as

$$\begin{cases} a \rightarrow a \\ b \rightarrow b \\ c \rightarrow c(\psi^q(a))^t. \end{cases}$$

Let $w = c(\psi^q(a))^t \in \mathbb{F}_3$ and let M be the length of w in the basis $\langle a, b, c \rangle$. Let $x \in \text{CV}_n$ be a rose where the edges are labelled a, b and w and

$$\ell_x(a) = \ell_x(b) = \frac{1}{M+2} \quad \text{and} \quad \ell_x(w) = \frac{M}{M+2}$$

and let $y \in \text{CV}_n$ be a rose where the edges are labelled a, b and c and have length $\frac{1}{3}$ each. Note that the length ratio of a, b and w from x to y are identical,

$$\frac{\ell_y(w)}{\ell_x(w)} = \frac{M/3}{M/(M+2)} = \frac{M+2}{3} = \frac{1/3}{1/(M+2)} = \frac{\ell_y(a)}{\ell_x(a)} = \frac{\ell_y(b)}{\ell_x(b)}.$$

In particular, we have

$$d_{CV_n(x,y)} = \log \frac{M+2}{3}.$$

In fact, it follows from [29] that there is a unique geodesic $[x, y]_{CV_n}$ in CV_n connecting x to y and it folds along the unique illegal turn. Namely, it folds the edge labelled w around the free factor \mathcal{A} and if $p: [0, m] \rightarrow \text{Out}(\mathbb{F}_n)$ is the shadow of $[x, y]_{CV_n}$ in $\text{Out}(\mathbb{F}_n)$, then the projection of p to $\mathcal{A} = \langle a, b \rangle$ backtracks. Hence, as was seen in the proof of Theorem 4.1, for any K and C , we can choose q and t large enough so that p is not a (K, C) -quasi-geodesics. \square

We can also modify the example in Theorem 4.1 to prove part (ii) of Theorem D. For brevity, we do not define greedy folding paths here. They are used in [5] in an essential way to prove the hyperbolicity of the free-factor graph. What we need is that if $\text{stretch}(x, y) = x$ then there is a greedy folding path connecting x to y .

THEOREM 5.2. *There are points x and y in the thick part of CV_n , such that they are connected by a greedy folding path whose shadow in $\text{Out}(\mathbb{F}_n)$ is not a quasi-geodesic.*

Proof. Let $\langle a, b, c \rangle$ be a generating set for \mathbb{F}_3 , let R_0 be a rose with labels $\{a, b, c\}$. Let ψ and $\psi_{\mathcal{A}}$ be as before. Let $y \in CV_n$ be the rose R_0 with edge labels a, b and c and edge lengths $\frac{1}{3}$.

Let \mathfrak{g} be the axis of $\psi_{\mathcal{A}}$ in $\mathcal{F}(\mathcal{A})$, the free factor graph associated to \mathcal{A} . Also let $\alpha_1, \alpha_2, \dots, \alpha_k \in \mathcal{A}$ be primitive loops in \mathcal{A} that, considered as vertices in $\mathcal{F}(\mathcal{A})$, are distance D or further from \mathfrak{g} for a large constant D .[†] For positive integers n_1, n_2, \dots, n_k , let ϕ be the following automorphism:

$$\begin{cases} a \longrightarrow \psi^n(a) \\ b \longrightarrow \psi^n(b) \\ c \longrightarrow c \alpha_1^{n_1} \cdots \alpha_k^{n_k}. \end{cases}$$

We observe that $|\psi^n(a)|_y$ and $|\psi^n(b)|_y$, as a function n , grow at a fixed exponential rate that is less than 3. Therefore for any given loops α_i and powers n_i , there is a power n so that

$$\max(|\psi^n(a)|_y, |\psi^n(b)|_y) \leq |c \alpha_1^{n_1} \cdots \alpha_k^{n_k}|_y \leq 3 \min(|\psi^n(a)|_y, |\psi^n(b)|_y). \tag{14}$$

Let x be a rose with edge labels, $\psi^n(a)$, $\psi^n(b)$ and $c \alpha_1^{n_1} \cdots \alpha_k^{n_k}$ (note that these form a basis for \mathbb{F}_3) and edge lengths

$$\frac{|\psi^n(a)|_y}{T}, \quad \frac{|\psi^n(b)|_y}{T} \quad \text{and} \quad \frac{|c \alpha_1^{n_1} \cdots \alpha_k^{n_k}|_y}{T},$$

where

$$T = |\psi^n(a)|_y + |\psi^n(b)|_y + |c \alpha_1^{n_1} \cdots \alpha_k^{n_k}|_y.$$

Then $\text{stretch}(x, y) = x$ and equation (14) implies that every edge length in x is larger than $\frac{1}{7}$ (x is $\frac{1}{7}$ -thick part of CV_n).

Let ϕ_i be the shortest automorphism such that $\phi_i(\alpha_i) = a$ and let $r_i = \|\phi_i\|$. It follows from construction that

$$\|\phi\| \leq 3n + \sum_{i=1}^k (2r_i + n_i). \tag{15}$$

[†]As we shall see, it is enough to have one such loop, however choosing many loops showcases how different the shadow of a quasi-geodesic in $\text{Out}(\mathbb{F}_n)$ to $\mathcal{F}(\mathcal{A})$ could be from being a quasi-geodesic.

Let ℓ_i be the length of α_i in the $\langle a, b \rangle$ basis and let λ be the stretch factor of ϕ . To make equation (14) hold, we need

$$\lambda^n \sim 1 + \sum_{i=1}^k n_i \ell_i.$$

Letting $n_{\max} = \max_i n_i$ we have, for constants c_1 and c_2 depending on k , r_i and ℓ_i , that

$$n \leq c_1 \log n_{\max}, \quad \text{and} \quad \|\phi\| \leq (k+1)n_{\max} + c_2.$$

Now let $p: [0, m] \rightarrow \text{Out}(\mathbb{F}_n)$ be the shadow of $[x, y]_{\text{gf}}$ to $\text{Out}(\mathbb{F}_n)$. Bestvina–Feighn[5] showed that the projection of $[x, y]_{\text{gf}}$ to $\mathcal{F}(\mathcal{A})$ is a quasi-geodesic. Hence the projection of $p[0, m]$ to $\mathcal{F}(\mathcal{A})$ stays near \mathfrak{g} and thus remains far from every α_i . And, again by [5, Lemma 3.3], this implies that the combinatorial length of α_i at any point along $[x, y]_{\text{gf}}$ is large, say larger than some constant L depending linearly on D . By Theorem 3.4,

$$m \succ \max_i n_i \log L.$$

But, since p is a quasi-geodesic, $m \prec \|\phi\|$. For large enough L , the above inequality contradicts equation (15). This implies p cannot be a quasi-geodesic. \square

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