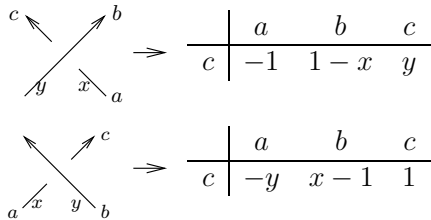


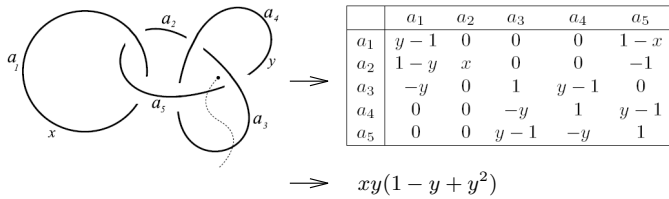
# The Penultimate Alexander Invariant

## A Definition of the MVA (From [Ar])

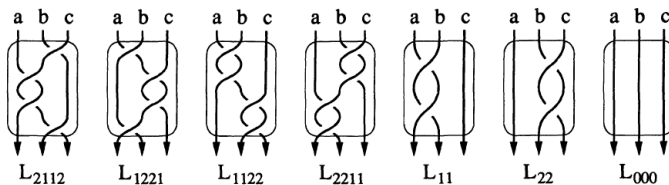
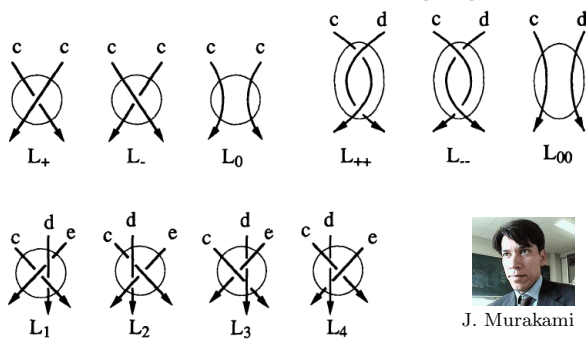


Joint with  
Jana Archibald

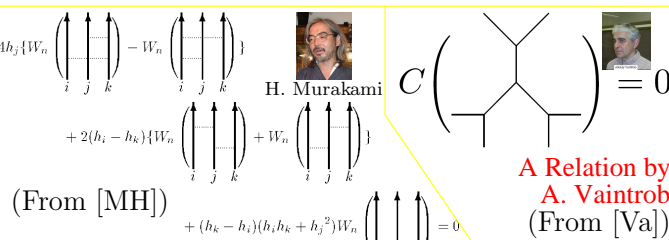
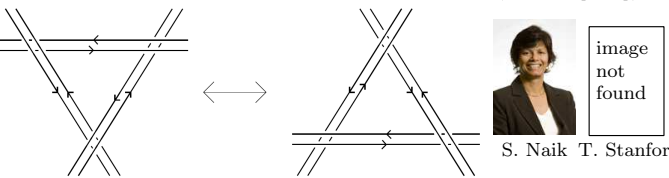
$$A = \frac{(-1)^{i+j} \det(M_i^j)}{w_i(t_i-1)} \prod_k t_k^{\frac{\text{rot}(k)-\mu(k)}{2}}$$



## Relations by J. Murakami (From [MJ])



## The Naik-Stanford Double Delta Relation (From [NS])



## A Relation by H. Murakami

There's Lots More!

"God created the knots,  
all else in topology  
is the work of mortals"  
Leopold Kronecker (paraphrased)



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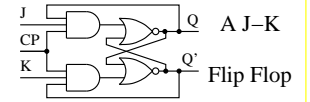
This handout and further links are at  
<http://www.math.toronto.edu/~drorbn/Talks/Sandbjerg-0810/>

Our Goal: Prove all these relations are consistent, at maximal confidence and minimal brain utilization.

⇒ We need an "Alexander Invariant" for arbitrary tangles, easy to define and compute and well-behaved under tangle compositions; better, "virtual tangles".

## Circuit Algebras

- \* Have "circuits" with "ends",
- \* Can be wired arbitrarily.
- \* May have "relations" – de-Morgan, etc.



Example  $VT = CA \langle \times, \times \rangle / R23 = PA \langle \times, \times, \times \rangle / R23, VR123, MR3$

Reminders from linear algebra. If  $X$  is a (finite) set,

$$\Lambda^k(X) := \langle k\text{-tuples in } X, \text{ modulo anti-symmetry} \rangle$$

$$\Lambda^{\text{top}}(X) := \langle |X|\text{-tuples in } X, \text{ modulo anti-symmetry} \rangle$$

$$\Lambda^{1/2}(X) := \langle (|X|/2)\text{-tuples in } X, \text{ modulo anti-symmetry} \rangle.$$

If  $Y \subset X^m$ , the "interior multiplication"  $i_Y : \Lambda^k(X) \rightarrow \Lambda^{k-m}(X)$  is anti-symmetric in  $Y$ .

**Definition.** An "Alexander half density with input strands  $X^{\text{in}}$  and output strands  $X^{\text{out}}$ " is an element of

$$\text{AHD}(X^{\text{in}}, X^{\text{out}}) := \Lambda^{\text{top}}(X^{\text{out}}) \otimes \Lambda^{1/2}(X^{\text{in}} \cup X^{\text{out}}).$$

Often we extend the coefficients to some polynomial ring without warning.

**Definition.** If  $\alpha_i \otimes p_i \in \text{AHD}(X_i^{\text{in}}, X_i^{\text{out}})$  (for  $i = 1, 2$ ), and  $G = (X_1^{\text{in}} \cup X_2^{\text{in}}) \cap (X_1^{\text{out}} \cup X_2^{\text{out}})$  is the set of "gluable legs", the "gluing" in  $\text{AHD}(X_1^{\text{in}} \cup X_2^{\text{in}} - G, X_1^{\text{out}} \cup X_2^{\text{out}} - G)$  is

$$i_G(\alpha_1 \wedge \alpha_2) \otimes i_G(p_1 \wedge p_2).$$

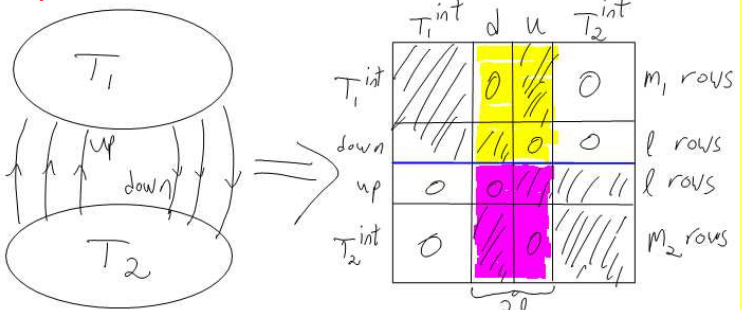
**Claim.** This makes AHD a circuit algebra.

**Definition.** The "Penultimate Alexander Invariant" is defined using

$$pA : \times_{l,i}^{k,j} \mapsto (j \wedge k) \otimes \begin{pmatrix} l \wedge i + (t_i - 1)l \wedge j - t_{il} \wedge k \\ + i \wedge j + t_{ij} \wedge k \end{pmatrix}$$

$$pA : \times_{i,j}^{l,k} \mapsto (k \wedge l) \otimes \begin{pmatrix} t_{ji} \wedge j - t_{ji} \wedge l + j \wedge k \\ + (t_i - 1)j \wedge l + k \wedge l \end{pmatrix}$$

## Why Works?



Every "rook arrangement" in the above picture must have exactly  $l$  rooks in the yellow zone and  $l$  rooks in the purple zone. So for  $T_1$  we only care about the minors in which exactly  $l$  of the  $2l$  middle columns are dropped, and the rest is signs...

**Weaknesses.** Exponential, no understanding of cablings, no obvious "meaning". The ultimate Alexander invariant should address all that...

**Challenge.** Can you categorify this?

```

1      (* WP: Wedge Product *)
2  WSort[expr_] := Expand[expr /. w_W -> Signature[w]*Sort[w]];
3  WP[0, _] = WP[_ , 0] = 0;
4  WP[a_, b_] := WSort[Distribute[a ** b] /.
5      (c1_. * w1_W) ** (c2_. * w2_W) -> c1 c2 Join[w1, w2]];
6
7      (* IM: Interior Multiplication *)
8  IM[{}, expr_] := expr;
9  IM[i_, w_W] := If[FreeQ[w, i], 0,
10     -(-1)^Position[w, i][[1,1]]*DeleteCases[w, i]];
11 IM[{is___, i_}, w_W] := IM[{is}, IM[i, w]];
12 IM[is_List, expr_] := expr /. w_W -> IM[is, w]
13
14     (* pA on Crossings *)
15 pA[Xp[i_,j_,k_,l_]] := AHD[(t[i]==t[k])(t[j]==t[l]), {i,l}, W[j,k],
16     W[l,i] + (t[i]-1)W[l,j] - t[l]W[l,k] + W[i,j] + t[l]W[j,k]];
17 pA[Xm[i_,j_,k_,l_]] := AHD[(t[i]==t[k])(t[j]==t[l]), {i,j}, W[k,l],
18     t[j]W[i,j] - t[j]W[i,l] + W[j,k] + (t[i]-1)W[j,l] + W[k,l]]
19
20     (* Variable Equivalences *)
21 ReductionRules[Times[]] = {};
22 ReductionRules[Equal[a_, b_]] := (# -> a)& /@ {b};
23 ReductionRules[eqs_Times] := Join @@ (ReductionRules /@ List@@eqs)
24
25     (* AHD: Alexander Half Densities *)
26 AHD[eqs_, is_, -os_, p_] := AHD[eqs, is, os, Expand[-p]];
27 AHD /: Reduce[AHD[eqs_, is_, os_, p_]] :=
28     AHD[eqs, Sort[is], WSort[os], WSort[p /. ReductionRules[eqs]]];
29 AHD /: AHD[eqs1_, is1_, os1_, p1_] AHD[eqs2_, is2_, os2_, p2_] := Module[
30     {glued = Intersection[Union[is1, is2], List@@Union[os1, os2]]},
31     Reduce[AHD[
32         eqs1*eqs2 /. eq1_Equal*eq2_Equal /;
33         Intersection[List@@eq1, List@@eq2] != {} -> Union[eq1, eq2],
34         Complement[Union[is1, is2], glued],
35         IM[glued, WP[os1, os2]],
36         IM[glued, WP[p1, p2]]
37     ] ] ]
38
39     (* pA on Circuit Diagrams *)
40 pA[cd_CircuitDiagram, eqs___] := pA[cd, {}, AHD[Times[eqs], {}, W[], W[]]];
41 pA[cd_CircuitDiagram, done_, ahd_AHD] := Module[
42     {pos = First[Ordering[Length[Complement[List @@ #, done]] & /@ cd]}],
43     pA[Delete[cd, pos], Union[done, List @@ cd[[pos]], ahd*pA[cd[[pos]]]]
44 ];
45 pA[CircuitDiagram[], _, ahd_AHD] := ahd

```

```

In[10]= Timing[res4 = pA[#, (t[1] == t[6]) (t[11] == t[16]) (t[21] == t[26])] & /@ {
CircuitDiagram[
  Xp[19, 1, 20, 2], Xp[11, 3, 12, 2], Xp[3, 30, 4, 29], Xp[4, 21, 5, 22],
  Xp[6, 23, 7, 22], Xp[7, 28, 8, 29], Xp[12, 8, 13, 9], Xp[18, 10, 19, 9],
  Xp[27, 13, 28, 14], Xp[23, 15, 24, 14], Xp[24, 16, 25, 17], Xp[26, 18, 27, 17]
],
CircuitDiagram[
  Xp[1, 28, 2, 27], Xp[2, 23, 3, 24], Xp[17, 3, 18, 4], Xp[13, 5, 14, 4],
  Xp[14, 6, 15, 7], Xp[16, 8, 17, 7], Xp[8, 25, 9, 24], Xp[9, 26, 10, 27],
  Xp[29, 11, 30, 12], Xp[21, 13, 22, 12], Xp[22, 18, 23, 19], Xp[28, 20, 29, 19]
]
]
]

```

A very large output was generated. Here is a sample of it:

```

Out[10]= {9.86, {AHD[(t[1] == t[2] == t[3] == t[4] == t[5] == t[6] == t[7] == t[8] == t[9] == t[10])
(t[11] == t[12] == t[13] == t[14] == t[15] == t[16] == t[17] == t[18] == t[19] == t[20])
(t[21] == t[22] == t[23] == t[24] == t[25] == t[26] == t[27] == t[28] == t[29] == t[30]),
{1, 6, 11, 16, 21, 26}, <<1>>,
-t[1]^2 t[11]^2 t[21]^2 W[1, 5, 6, 11, 15, 21] + <<2574>>, <<1>>}}

```

Show Less Show More Show Full Output Set Size Limit...

In[11]= Equal @@ (Last /@ res4)

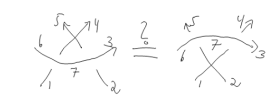
Out[11]= True

(The program also prints "False" when appropriate, and computes Alexander polynomials.)

More at <http://www.math.toronto.edu/~drorbn/Sandbjerg-0810/pA.nb>

$i_1 \wedge i_2 \wedge \dots$ . To sort it we Sort its arguments and multiply by the Signature of the permutation used. **3.** The wedge product of 0 with anything is 0. **4-5.** The wedge product of two things involves applying the Distributive law, Joining all pairs of W's, and WSorting the result. **8.** Inner multiplying by an empty list of indices does nothing. **9-10.** Inner multiplying a single index yields 0 if that index is not present, otherwise it's a sign and the index is deleted. **11-12.** Afterwards it's simple recursion. **15-18.** For the crossings Xp and Xm it is straightforward to determine the incoming strands, the outgoing ones, and the variable equivalences. The associated half-densities are just as in the formulas. **21-23.** The technicalities of imposing variable equivalences are annoying. **26.** That's all we need from the definition of a tensor product. **27-28.** Straightforward simplifications. **29.** The (circuit algebra) product of two Alexander Half Densities: **30.** The glued strands are the intersection of the ins and the outs. **32-33.** Merging the variable equivalences is tricky but natural. **34-35.** Removing the glued strands from the ins and outs. **36 The Key Point.** The wedge product of the half-densities, inner with the glued strands. **40-45.** A quick implementation of a "thin scanning" algorithm for multiple products. The key line is **42**, where we select the next crossing we multiply in to be the crossing with the fewest "loose strands".

Overcrossings Commute



Hence "w-knots"

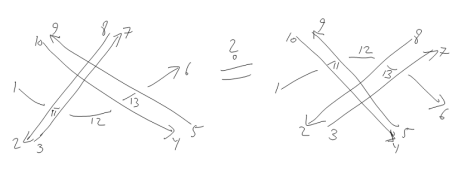
```

In[15]= Equal[
  pA[CircuitDiagram[Xp[1, 7, 4, 6], Xp[2, 3, 5, 7]]],
  pA[CircuitDiagram[Xp[2, 7, 5, 6], Xp[1, 3, 4, 7]]]]

```

Out[15]= True

Commutators Commute



Question.

Does this specify the Alexander polynomial?

```

In[5]= Equal[
  pA[CircuitDiagram[Xp[1, 2, 11, 8], Xm[11, 3, 12, 7],
  Xp[12, 4, 13, 10], Xm[13, 5, 6, 9]], t[2] == t[3], t[4] == t[5]],
  pA[CircuitDiagram[Xp[1, 4, 11, 10], Xm[11, 5, 12, 9],
  Xp[12, 2, 13, 8], Xm[13, 3, 6, 7]], t[2] == t[3], t[4] == t[5]]]

```

Out[5]= True

## References

[Ar] J. Archibald, *The Weight System of the Multivariable Alexander Polynomial*, arXiv:0710.4885.

[MH] H. Murakami, *A Weight System Derived from the Multivariable Conway Potential Function*, Jour. of the London Math. Soc. **59** (1999) 698–714, arXiv:math/9903108.

[MJ] J. Murakami, *A State Model for the Multi-Variable Alexander Polynomial*, Pac. Jour. of Math. **157-1** (1993) 109–135.

[NS] S. Naik and T. Stanford, *A Move on Diagrams that Generates S-Equivalence of Knots*, Jour. of Knot Theory and its Ramifications **12-5** (2003) 717–724, arXiv:math/9911005.

[Va] A. Vaintrob, *Melvin-Morton Conjecture and Primitive Feynman Diagrams*, Inter. J. Math. **8** (1997) 537–553, arXiv:q-alg/9605028.