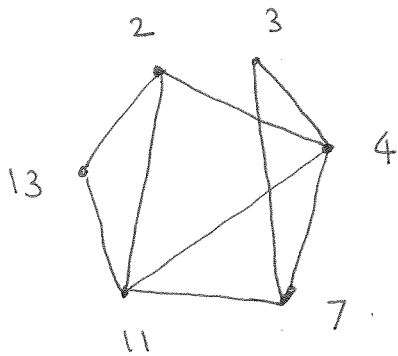


1.3



1.12

(a)  $x u v r u v y$

(b)  $v s t v w$

(c) No such path: neighbors of  $r = \{u, s\}$   
and  $z$  is not a neighbor of  $u$  or  $s$ .

(d) No such path.

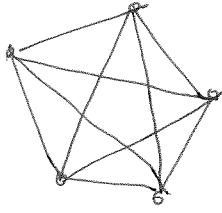
(e)  $x u v t$  ( $d(x, t) = 3$ )

(f)  $x u r v y x u r v y x$

(g)  $x u r s t w z y x$

(h)  $r v y z$ ,  $\text{diam}(G) = 3$

1.15.



All distances have to be 1, because if some distance  $d(x, y) \geq 2$ , then the oddness condition means that  $d(x, y) \geq 3$  and  $d(x, y)$  is odd.

But if  $x, a_1, a_2, \dots, a_{l-1}, y$  is a path of length  $d(x, y)$ , then

$$d(a_1, y) = d(x, y) - 1 \text{ and is } \text{even} \text{ and } \geq 2 \text{ (and hence } d(a_1, y) \neq 0).$$

This is a contradiction. ~~⊗~~

⊗

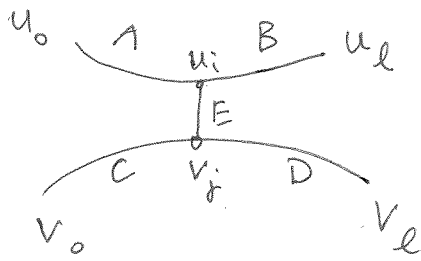
So the graph has to be  $K_5$ .

1.17. (a).

Suppose  $u_0, u_1, \dots, u_l$  and  $v_0, v_1, \dots, v_l$  are longest paths in  $G$  with no ~~⊗~~ common vertices.

Let  $u_i, v_j$  be ~~⊗~~ a pair of vertices that minimizes  $d(u_i, v_j)$   $i \in \{0, 1, \dots, l\}$   $j \in \{0, 1, \dots, l\}$ .

Then  $d(u_i, v_j) > 0$  by assumption.



Let  $A, B, C, D, E$  be the paths shown.

(so  $A$  is the path ~~from~~  $u_0, u_1, \dots, u_i$

$B$  is  $u_i, u_{i+1}, \dots, u_l$

and similarly for  $C, D$ .)

$E$  is the <sup>shortest</sup> path from  $u_i$  to  $v_j$ )

Note that  $\text{length}(E) = d(u_i, v_j) > 0$ .

~~QED~~

Consider the path

(longer one of  $A, B$ )

followed by

$E$

followed by

(longer one of  $C, D$ ).

The total length of this path is

$(\geq l/2) + (\geq 1) + (\geq l/2)$  which is

$\geq l+1$ , contradicting the hypothesis.

Let  $N_i = \#$  vertices with degree  $i$

2-3.  
12 vertices  
31 edges.

Handshake thm:

$$4 \cdot N_4 + 6 \cdot N_6 = \textcircled{2} \textcircled{3} 2 \cdot 31$$

$$N_4 + N_6 = 12$$

Solve:

$$N_4 = 5, \quad N_6 = 7$$

2.5.  
25 vertices.  
62 edges.

$$N_4 = 2$$

$$N_6 = 11.$$

$$N_3, N_5 ?$$

$$N_3 + N_4 + N_5 + N_6 \textcircled{=} \textcircled{2} \textcircled{3} \textcircled{4} 25$$

$$3N_3 + 4N_4 + 5N_5 + 6N_6 = 2 \cdot 62$$

Solve --

2.7.  
Each edge has one endpoint in  $U$  and one endpoint in  $V$ .

⊗ So  $\deg u$  counts the number of edges incident on  $u$ .  $\sum_{u \in U} (\deg u) = \text{total \# of edges} = m$ .  
Similarly for  $V$ .