

FAMILY NAME _____

GIVEN NAME(S) _____

STUDENT NUMBER _____

SIGNATURE _____

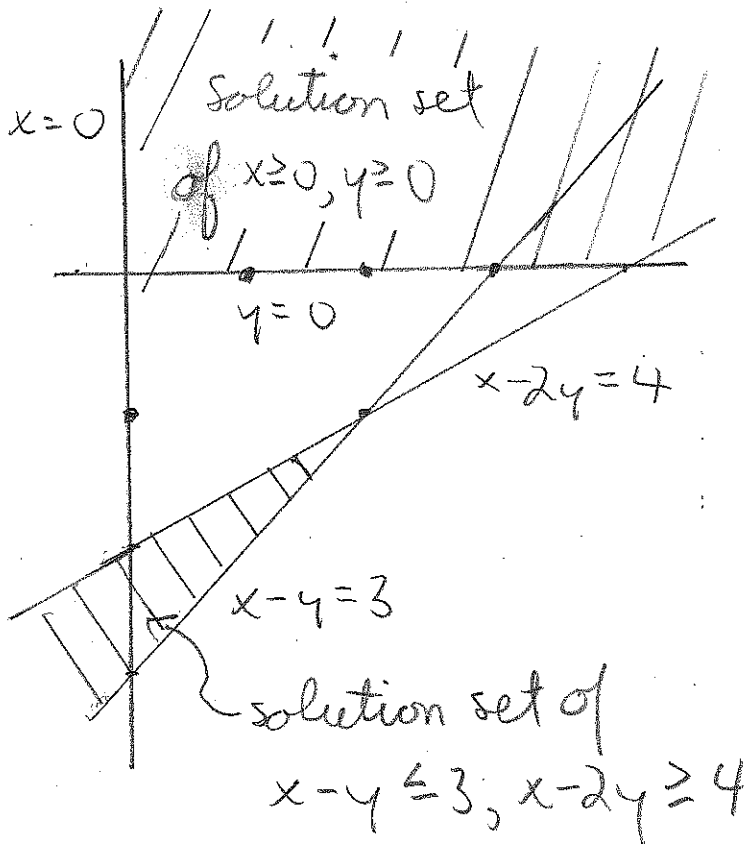
Instructions: No calculators or other aids allowed.

This test has 3 questions whose values are given immediately after the question numbers. Total marks = 40.

Write solutions in the spaces provided, using the backs of the pages if necessary. (Suggestion: If you have to continue a question, you may use the back of the **previous** page.) Aspects of any question which are indicated in **boldface** will be regarded as crucial during grading. **Show your work.**

The duration of this test is 50 minutes.

1. (13 marks) **Solve** the following problem **graphically**: Maximize $z = 5x + 6y$ subject to the constraints $x - y \leq 3$, $x - 2y \geq 4$, $x \geq 0, y \geq 0$.



The feasible region of the problem is the intersection of the two shaded regions. Since they do not intersect, the problem is infeasible.

2.(a) (4 marks) Which points in \mathbb{R}^n belong to the **line segment** having endpoints $x_1 \in \mathbb{R}^n$ and $x_2 \in \mathbb{R}^n$?

One correct answer is: points of the form $(1-\lambda)x_1 + \lambda x_2$ where $0 \leq \lambda \leq 1$

A second correct answer is: points of the form $c_1 x_1 + c_2 x_2$ where $c_1 + c_2 = 1$, $c_1 \geq 0$, $c_2 \geq 0$

2.(b) (5 marks) Define the term convex set (in \mathbb{R}^n).

One correct answer is: S is convex provided, for each x_1 and $x_2 \in S$, the line segment joining x_1 and x_2 lies in S .

A second correct answer is: S is convex provided, for each $x_1, x_2 \in S$ and $\lambda \in [0, 1]$, $(1-\lambda)x_1 + \lambda x_2 \in S$.

2.(c) (5 marks) Let $S = \left\{ \begin{bmatrix} x \\ y \end{bmatrix} \in \mathbb{R}^2 \text{ s.t. } -x + 2y \leq 2 \text{ and } x - y \leq 0 \right\}$. Prove that $\begin{bmatrix} 0 \\ 1 \end{bmatrix}$

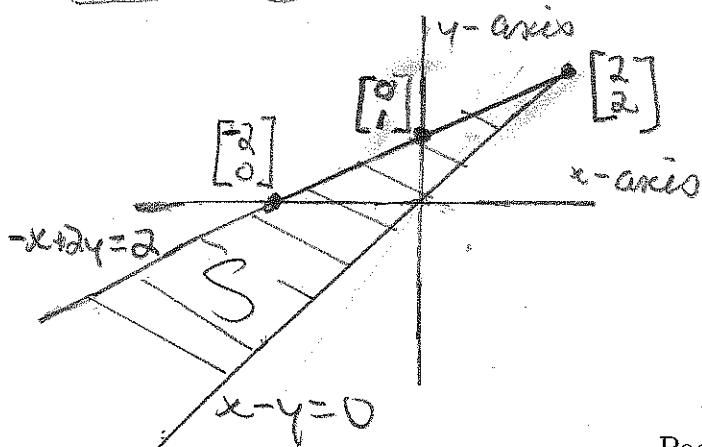
is not an extreme point of S .

$$\begin{bmatrix} -2 \\ 0 \end{bmatrix} \in S, \text{ since } -(-2) + 2 \cdot 0 \leq 2 \text{ and } -2 - 0 \leq 0$$

$$\begin{bmatrix} 2 \\ 2 \end{bmatrix} \in S, \text{ since } -2 + 2 \cdot 2 \leq 2 \text{ and } 2 - 2 \leq 0$$

$$\text{Also, } \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} -2 \\ 0 \end{bmatrix} + \frac{1}{2} \begin{bmatrix} 2 \\ 2 \end{bmatrix}, \text{ yet } \begin{bmatrix} 0 \\ 1 \end{bmatrix} \neq \begin{bmatrix} -2 \\ 0 \end{bmatrix} \text{ and } \begin{bmatrix} 0 \\ 1 \end{bmatrix} \neq \begin{bmatrix} 2 \\ 2 \end{bmatrix}$$

Q.E.D.



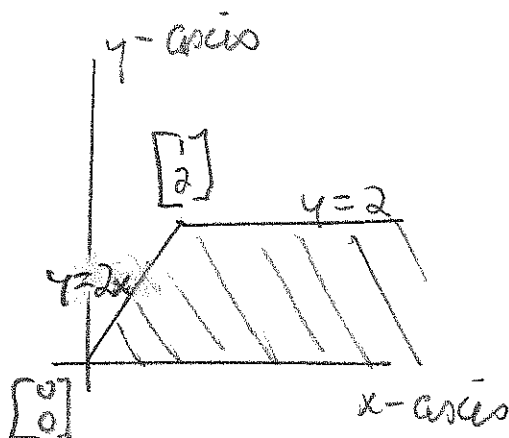
This diagram is not part of the proof completed above. It was used to motivate the proof.

3. (13 marks) Write **one** linear programming problem, \mathcal{P} , which satisfies **all** of the following:

- (1) \mathcal{P} has 2 decision variables, x and y .
- (2) \mathcal{P} is in **standard form**.
- (3) The feasible region of \mathcal{P} has $\begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$ and $\begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$ as **extreme points**.
- (4) $\begin{bmatrix} 0 \\ 0 \end{bmatrix}$ and $\begin{bmatrix} 1 \\ 2 \end{bmatrix}$ are the **only** extreme points of the feasible region of \mathcal{P} .
- (5) \mathcal{P} is **unbounded**.

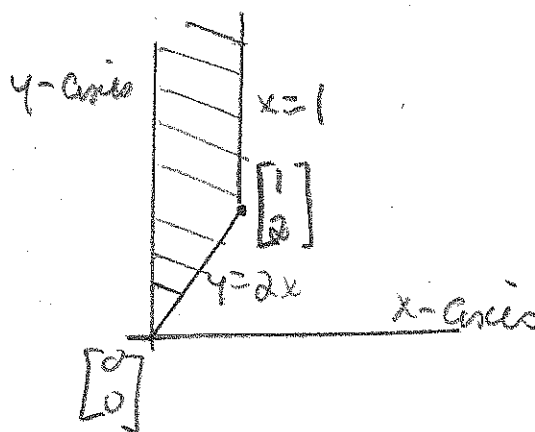
One solution is:

$$\begin{aligned} \text{Maximize } z &= x \text{ s.t.} \\ -2x + y &\leq 0 \\ y &\leq 2 \\ x \geq 0, y &\geq 0 \end{aligned}$$



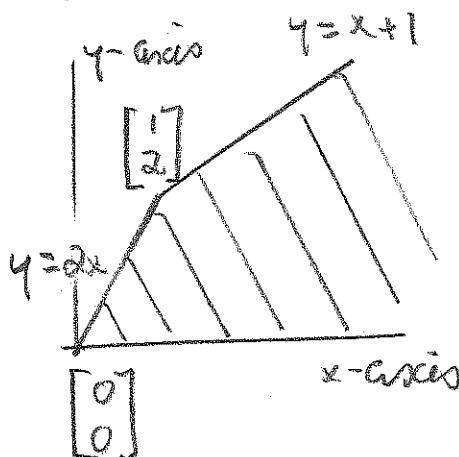
A second solution is:

$$\begin{aligned} \text{Maximize } z &= y \text{ s.t.} \\ 2x - y &\leq 0 \\ x &\leq 1 \\ x \geq 0, y &\geq 0 \end{aligned}$$



A third solution is:

$$\begin{aligned} \text{Maximize } z &= x \text{ s.t.} \\ -2x + y &\leq 0 \\ -x + y &\leq 1 \\ x \geq 0, y &\geq 0 \end{aligned}$$



(Questions 2. (c) and 3 each have infinitely many solutions.)