Problem 1. [6 points] Give a combinatorial proof for the formula, where $x, y \in \mathbb{N}$.

$$\sum_{k=0}^{n} \binom{x}{k} \binom{y}{n-k} = \binom{x+y}{n}$$

Solution:

- (2 points) The right hand side $\binom{x+y}{n}$ naturally counts *n*-element subsets of the set [x+y]. The left hand side can be interpreted as counting pairs of subsets (S_1, S_2) with $S_1 \subseteq [x], S_2 \subseteq [y], |S_1| + |S_2| = n$. We will set up a bijection between the two sides.
- (2 points) Given an *n*-subset $S \subseteq [x + y]$, the intersection $S \cap [x]$ is a subset of [x]. Let $k = |S \cap [x]|$. Clearly $0 \le k \le x$. The complementary subset $S \setminus [x]$ is a subset of $\{x + 1, x + 2, \dots, x + y\}$ with n k many elements. We may identify $S \setminus [x]$ with a subset of [y] by subtracting x from each element (subtracting a constant from every element of a set is clearly a bijection). So we defined a map f sending an n-subset S of [x + y] to a pair of subsets (S_1, S_2) with $S_1 \subseteq [x], S_2 \subseteq [y], |S_1| + |S_2| = n$.
- (2 points) We now show that f is a bijection. To show that f is surjective, note that given any pair of subsets (S_1, S_2) with the conditions above, we can form the set $S_1 \cup (S_2+x)$, where $S_2+x = \{j+x \mid j \in S_2\}$, and $S_1 \cup (S_2+x) \subseteq [x+y]$ with $|S_1 \cup (S_2+x)| = n$. We clearly have $f(S_1 \cup (S_2+x)) = (S_1, S_2)$. To show that f is injective, note that if f(A) = f(B) for n-subsets A, B of [x+y], then we must have $A \cap [x] = B \cap [x]$, and also $A \cap ([n] \setminus [x]) = B \cap ([n] \setminus [x])$. Hence

$$A = (A \cap [x]) \cup (A \cap ([n] \setminus [x])) = (B \cap [x]) \cup (B \cap ([n] \setminus [x])) = B$$

so f is injective.

Solution:

- (2 points) The right hand side $\binom{x+y}{n}$ naturally counts binary strings of length x + y with exactly n 1s. The left hand side can be interpreted as counting pairs of binary strings (S_1, S_2) of length x and y with a combined number of n 1s. We will set up a bijection between the two sides.
- (2 points) Given a binary string of length x + y with exactly n 1s, we can separate it into two strings, the initial x-string $S_{\leq x}$ and the terminal y-string $S_{>x}$. Clearly the toal number of 1s is unchanged by this, so this defines a map f from binary strings of length x + y with n 1s to pairs of binary strings with a combined number of n 1s.
- (2 points) We now show f is a bijection. We will do this by defining an inverse map g. Given a pair (S_1, S_2) of binary strings of lengths x and y with a combined total of n 1s, let $g(S_1, S_2) = S_1S_2$ (the concatenation of the two strings). This defines a binary string of length x + y with n 1s. Also, we have

$$g(f(S)) = g(S_{\le x}, S_{>x}) = S_{\le x}S_{>x}$$

and

$$f(g(S_1, S_2)) = f(S_1 S_2) = ((S_1 S_2)_{\le x}, (S_1 S_2)_{>x}) = (S_1, S_2),$$

where the last equality follows from the fact that $|S_1| = x$ and $|S_2| = y = (x + y) - x$. Since f has a two-sided inverse, it must be a bijection.

Problem 2. Find a recurrence relation for the number g(n) of ternary strings (that is, $\{0, 1, 2\}$ -strings) of length n that do not contain 102 as a substring.

Solution:

- (1 point) Any ternary string of length n can be obtained from a ternary string of length n-1 by appending a character (a 0,1, or 2) at the end.
- (3 points) Given a ternary string of length n 1 that does not contain 102 as a substring, we can get an *n*-string from it by appending a character at the end. The resulting string will not contain the substring 102, unless we append a 2 at the end of an n 1 string that ends with 10. There are g(n 3) many such ternary strings that do not contain 102 as a substring.
- (2 points) Therefore, the recurrence relation is

$$g(n) = 3g(n-1) - g(n-3).$$