

# Lecture 11: Restricted Intersections of Sets

Combinatorial Methods (Winter 2023)  
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Fix some positive integers  $n, a, b$ . Consider families of sets  $\mathcal{Y} \subseteq \binom{[n]}{a}$  such that for all  $u, v \in \mathcal{Y}$ :

$$|u \cap v| \leq b$$

We are interested in finding out how big/small can the set  $\mathcal{Y}$  be, and in this lecture we present various approaches to this problem.

We write  $a = \alpha n$  and  $b = \beta n$  for some  $\alpha, \beta \in [0, 1]$ .

## 1 Conditions on $\beta$ for $|\mathcal{Y}|$ to be small

**Theorem 1.** *If  $\beta < 2\alpha - 1$ , then  $|\mathcal{Y}| \leq 1$ .*

*Proof.* Suppose otherwise. If  $u, v \in \mathcal{Y}$  are such that  $u \neq v$ , then:

$$\beta n = b \geq |u \cap v| = |u| + |v| - |u \cup v| \geq 2a - n = (2\alpha - 1)n$$

a contradiction. □

Conversely, if  $\beta > \alpha$ , then trivially  $\mathcal{Y}$  can be arbitrarily large (e.g. it is possible that  $\mathcal{Y} = \binom{[n]}{a}$ ). The growth rate is exponential: By Stirling's approximation for  $n!$ , we have that:

$$\binom{n}{\alpha n} \sim 2^{c_\alpha n} \cdot \text{poly}(n)$$

where  $\text{poly}(n)$  is some polynomial with variable  $n$ , and:

$$c_\alpha = H(\alpha) = \alpha \log \frac{1}{\alpha} + (1 - \alpha) \log \left( \frac{1}{1 - \alpha} \right)$$

## 2 Growth rates of $|\mathcal{Y}|$

We now consider the following questions:

1. For which  $\alpha, \beta$  can  $|\mathcal{Y}|$  be exponentially big?
2. For which  $\alpha, \beta$  must  $|\mathcal{Y}|$  be  $\leq \text{poly}(n)$ ?

## 2.1 Exponential growth rate

Fix some  $m$ , and write:

$$\mathcal{Y} = \{u_1, u_2, \dots, u_m\}$$

where each  $u_i$  is picked independently and uniformly from  $\binom{[n]}{a}$ . Since each number in  $[n]$  has a uniform probability of  $\frac{1}{n}$  to be in any of the  $v_i$ 's, we “expect” the “typical intersection” size to be  $\frac{a^2}{n} = \alpha^2 \cdot n$ . To see this: Instead of taking  $u \in \binom{[n]}{a}$ , we take each element of  $[n]$  into  $u$  with probability  $\frac{a}{n}$ . Thus, for a specific  $i \in [n]$  and two fixed sets in  $\binom{[n]}{a}$ , there is  $(\frac{a}{n})^2$  probability of  $i$  belonging in both sets. Since there are  $n$  many such  $i$ 's, the expected size of intersection is  $n \cdot (\frac{a}{n})^2$ .

This also tells us the following:

**Theorem 2.** *If  $\beta > \alpha^2$ , then there exists exponentially large  $\mathcal{Y}$ .*

Intuitively, this is because if we take a large family of random sets, then by the above reasoning most of the families should have intersection at most  $\alpha^2 n < b$ , which satisfies the requirement  $|u \cap v| \leq b$ .

*Proof.* Pick  $b$  such that  $a^2 < p < b$ . For each  $i \in [m]$ , where  $m$  is to be determined, let  $u_i$  be a set where each element is taken independently with probability  $p$ . Then:

$$\mathbb{E}[|u_i|] = \mathbb{E}\left[\sum_{k=1}^n \mathbb{1}_{k \in u_i}\right] = \sum_{k=1}^n \mathbb{E}[\mathbb{1}_{k \in u_i}] = pn$$

By Chernoff's bound, we have that for all  $i$  and  $\varepsilon > 0$ :

$$\Pr[||u_i| - pn| > \varepsilon n] \leq e^{-\frac{\varepsilon^2 n}{4}}$$

Choose  $\varepsilon := p - \alpha$ . Then:

$$\Pr[|u_i| < \alpha n] \leq \Pr[||u_i| - pn| > \varepsilon n] \leq e^{-\frac{\varepsilon^2 n}{4}}$$

Therefore:

$$\Pr[\exists i \in [m] \text{ such that } |u_i| < \alpha n] = \sum_{i=1}^m \Pr[|u_i| < \alpha n] \leq m e^{-\frac{\varepsilon^2 n}{4}}$$

Now fix  $i, j \in [m]$  with  $i \neq j$ . Then:

$$|v_i \cap v_j| = \mathbb{E}\left[\sum_{k=1}^n \mathbb{1}_{k \in u_i} \cdot \mathbb{1}_{k \in u_j}\right] = \sum_{k=1}^n \Pr[k \in u_i \wedge k \in u_j] = p^2 n$$

Let  $X_k := \mathbb{1}_{k \in u_i} \cdot \mathbb{1}_{k \in u_j}$  (so  $\Pr[X_k = 1] = p^2$ ). Then, by Chernoff's bound again:

$$\Pr\left[\left|\sum X_k - p^2 n\right| > \varepsilon^2 n\right] \leq e^{-\frac{\varepsilon^2 n}{4}} \tag{1}$$

Choose  $\epsilon = \beta - p^2$  here, and we have that:

$$\Pr[|u_i \cap u_j| > \beta n] \leq e^{-\frac{\epsilon^2 n}{4}}$$

Therefore:

$$\Pr[\exists i, j \in [m], i \neq j \text{ such that } |u_i \cap u_j| > \beta n] \leq \binom{m}{2} e^{-\frac{\epsilon^2 n}{4}} \quad (2)$$

Thus, if  $m$  is chosen such that  $me^{-\frac{\epsilon^2 n}{4}} < 1$  (inequality (1)) and  $\binom{m}{2}e^{-\frac{\epsilon^2 n}{4}} < 1$  (inequality (2)), then with positive probability we obtain a family with size of intersections at most  $\beta n = b$ .  $\square$

## 2.2 Polynomial growth rate

**Theorem 3.** *If  $\beta < \alpha^2$ , then  $|\mathcal{Y}| \leq n + 1$ .*

*Proof.* We consider each  $u_i$  as a vector  $u_i \in \{0, 1\}^n \subseteq \mathbb{R}^n$ , such that for all  $i \neq j$ :

$$\langle u_i, u_i \rangle = \alpha n, \quad \langle u_i, u_j \rangle \leq \beta n$$

Write  $\mathcal{Y} = \{u_1, \dots, u_m\}$ .

**Lemma 4.** *If  $\tilde{u}_1, \dots, \tilde{u}_m \in \mathbb{R}^n \setminus \{0\}$  are such that  $\langle \tilde{u}_i, \tilde{u}_j \rangle < 0$  for all  $i \neq j$ , then  $m \leq n + 1$ .*

*Proof.* We induct on  $n$ . Assume WLOG that  $\tilde{u}_i$  are all unit vectors, and by rotating the space if necessary we also assume that  $\tilde{u}_1 = e_1$ . Then write out the vectors as follow:

$$\begin{aligned} \tilde{u}_1 &= [1 \quad 0 \quad \cdots \quad 0] \\ \tilde{u}_2 &= [c_2 \quad -\vec{r}_2 \quad -] \\ \tilde{u}_3 &= [c_3 \quad -\vec{r}_3 \quad -] \\ &\vdots \\ \tilde{u}_m &= [c_m \quad -\vec{r}_m \quad -] \end{aligned}$$

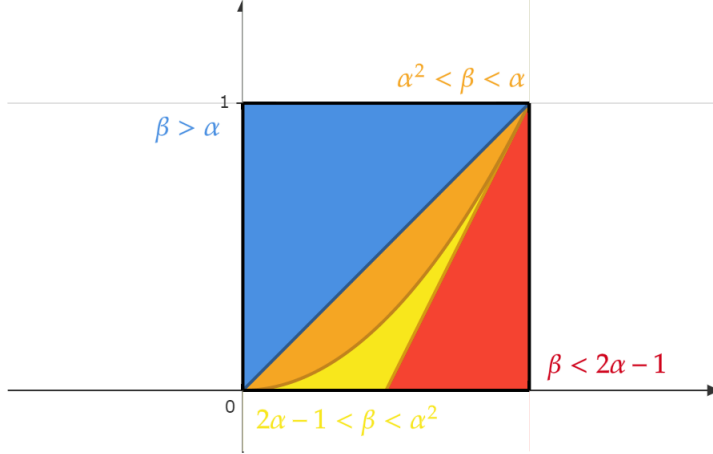
We have that  $c_2, \dots, c_m < 0$ , and  $\vec{r}_2, \dots, \vec{r}_m \in \mathbb{R}^{n-1}$ . This implies that  $\langle \vec{r}_i, \vec{r}_j \rangle < 0$  for all  $i \neq j$ . By induction hypothesis, we have that  $m - 1 \leq n$ , so  $m \leq n + 1$ .  $\blacksquare$

We now prove the theorem. Take  $w = \gamma \cdot \vec{1} \in \mathbb{R}^n$  (where  $\vec{1} = (1, \dots, 1)$ ), where  $\gamma$  is to be determined. Let  $\tilde{u}_i := u_i - w$ . Then for  $i \neq j$ :

$$\begin{aligned} \langle \tilde{u}_i, \tilde{u}_j \rangle &= \langle u_i - w, u_j - w \rangle \\ &= \langle u_i, u_j \rangle - \langle u_i, w \rangle - \langle u_j, w \rangle + \langle w, w \rangle \\ &\leq \beta n - \gamma \alpha n - \gamma \alpha n + n \gamma^2 \\ &= \gamma^2 - 2\alpha n \gamma + \beta n \\ &= n [(\gamma - \alpha)^2 + (\beta - \alpha^2)] \end{aligned}$$

Choose  $\gamma = \alpha$ , and we have that  $\langle \tilde{u}_i, \tilde{u}_j \rangle \leq n(\beta - \alpha^2) < 0$ . By the Lemma above, we have that  $m \leq n + 1$ , as desired.  $\square$

In summary, we have the following:



1. If  $\beta < \alpha$ , then the family can be arbitrarily large.
2. If  $\alpha^2 < \beta < \alpha$ , then by Theorem 2 the family can be exponentially large.
3. If  $2\alpha - 1 < \beta < \alpha$ , then by Theorem 3 the family can be polynomially large.
4. If  $\beta < \alpha$ , then by Theorem 1 the family has at most one set.

### 3 Bounding $|\mathcal{Y}|$ with using bipartite graphs

Again, let  $\mathcal{Y} = \{u_1, \dots, u_m\} \subseteq \binom{[n]}{a}$ . Consider a bipartite graph  $G = (\mathcal{Y} \sqcup [n], E)$ , such that for each  $i$  and  $k$ ,  $(u_i, k) \in E$  iff  $k \in u_i$ . The assertion that  $|u_i \cap u_j| \leq \beta n$  for all  $i \neq j$  is equivalent to saying that the graph  $G$  does not contain  $K_{2, \beta n + 1}$  as an induced subgraph. How can we use this property to bound  $m$ ?

We start by counting the number of subgraphs of the shape “>”. That is, an induced subgraph with vertices  $\{u_i, k, u_j\}$  for some  $u_i, u_j \in \mathcal{Y}$  and  $k \in [n]$ . We first note that:

$$\#(\text{“>”}) \leq \binom{m}{2} \cdot \beta n = \frac{m(m-1)\beta n}{2}$$

as there are  $\binom{m}{2}$  many pairs in  $\mathcal{Y}$ , and each pair has at most  $\beta n$  many elements in their intersection. On the other hand:

$$\#(\text{“>”}) = \sum_{k \in [n]} \binom{\deg(k)}{2} \geq \sum_{k \in [n]} \binom{\overline{\deg}}{2} = n \cdot \binom{\alpha m}{2} = \frac{n(\alpha m)(\alpha m - 1)}{2}$$

where  $\overline{\text{deg}}$  is the average degrees of nodes on the right. Combining both inequalities, we have that:

$$\begin{aligned} \frac{n(\alpha m)(\alpha m - 1)}{2} &\leq \frac{m(m - 1)\beta n}{2} \implies m(\alpha^2 - \beta) \leq \alpha - \beta \\ &\implies m \leq \frac{\alpha - \beta}{\alpha^2 - \beta} \end{aligned}$$