

# Handout on relations (projection, restriction, join, density)

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## 1 Relations

We consider relations over a fixed universe  $[n]$  ( $= \{1, \dots, n\}$ ). Rather than  $k$ -tuples and  $k$ -ary relations over  $[n]$  (i.e. elements and subsets of  $[n]^k$  where  $k \in \mathbb{N}$ ), we will speak of tuples and relations over  $[n]$  indexed by finite sets. Below, let  $V, W, S, T$  represent arbitrarily finite sets.

**Definition 1** ( $V$ -Tuples).

- A  $V$ -tuple is an element  $x \in [n]^V$ . For  $v \in V$ , the  $v$ -th coordinate of  $x$  is denoted  $x_v \in [n]$ .
- There is a unique  $\emptyset$ -ary tuple, which we denote by  $\circ$ . That is,  $[n]^\emptyset = \{\circ\}$ .
- For  $x \in [n]^V$  and  $S \subseteq V$ , we write  $x_S \in [n]^S$  for the restriction of  $x$  to coordinates in  $S$ .
- For  $y \in [n]^S$  and  $z \in [n]^{V \setminus S}$ , we write  $yz$  for the  $V$ -tuple  $x \in [n]^V$  with  $x_S = y$  and  $x_{V \setminus S} = z$ . (This concatenation operation is unordered, that is,  $yz = zy$ .)

**Definition 2** ( $V$ -ary Relations, Projection, Restriction).

- A  $V$ -ary relation is a set  $\mathcal{A} \subseteq [n]^V$  of  $V$ -tuples.
- For  $\mathcal{A} \subseteq [n]^V$  and  $S \subseteq V$ , the  $S$ -projection of  $\mathcal{A}$  is the  $S$ -ary relation

$$\text{proj}_S(\mathcal{A}) := \{x_S : x \in \mathcal{A}\}.$$

- For  $\mathcal{A} \subseteq [n]^V$  and  $S \subseteq V$  and  $z \in [n]^{V \setminus S}$ , the  $S$ -restriction of  $\mathcal{A}$  at  $z$  is the  $S$ -ary relation

$$\mathcal{A}|_S^z := \{y \in [n]^S : yz \in \mathcal{A}\}.$$

Note that  $\mathcal{A}|_S^z \subseteq \text{proj}_S(\mathcal{A})$ .

**Definition 3** (Join).

- The *join* of relations  $\mathcal{A} \subseteq [n]^V$  and  $\mathcal{B} \subseteq [n]^W$  is the  $V \cup W$ -ary relation defined by

$$\mathcal{A} \bowtie \mathcal{B} := \{x \in [n]^{V \cup W} : x_V \in \mathcal{A} \text{ and } x_W \in \mathcal{B}\}.$$

The join operation  $\bowtie$  is a hybrid of intersection  $\cap$  and product  $\times$ . If  $V = W$ , then we have  $\mathcal{A} \bowtie \mathcal{B} = \mathcal{A} \cap \mathcal{B}$ . On the opposite extreme, if  $V \cap W = \emptyset$ , then we have  $\mathcal{A} \bowtie \mathcal{B} = \mathcal{A} \times \mathcal{B}$  (i.e. the set  $\{xy : x \in \mathcal{A} \text{ and } y \in \mathcal{B}\}$ ). As an operation on relations, note that  $\bowtie$  is associative, commutative and idempotent (i.e.  $\mathcal{A} \bowtie (\mathcal{B} \bowtie \mathcal{C}) = (\mathcal{A} \bowtie \mathcal{B}) \bowtie \mathcal{C}$  and  $\mathcal{A} \bowtie \mathcal{B} = \mathcal{B} \bowtie \mathcal{A}$  and  $\mathcal{A} \bowtie \mathcal{A} = \mathcal{A}$ ). Also note that  $\emptyset$ -ary relation  $\{\circ\}$  is an identity under  $\bowtie$  (i.e.  $\mathcal{A} \bowtie \{\circ\} = \mathcal{A}$ ).

**Exercise 4.** Suppose  $\mathcal{A} \subseteq [n]^V$  and  $\mathcal{B} \subseteq [n]^W$  and  $\mathcal{C} \subseteq [n]^{V \cup W}$  such that  $\mathcal{C} \subseteq \mathcal{A} \bowtie \mathcal{B}$ .

- Show that  $\text{proj}_V(\mathcal{C}) \subseteq \mathcal{A}$ .
- For  $S \subseteq V$ , show that  $\text{proj}_S(\mathcal{C}) \subseteq \text{proj}_S(\mathcal{A})$ .

## 2 Density

**Definition 5** (Density, Projection Density, Restriction Density).

- The *density* of a relation  $\mathcal{A} \subseteq [n]^V$ , denoted  $\mu(\mathcal{A})$ , is defined by

$$\mu(\mathcal{A}) := \frac{|\mathcal{A}|}{n^{|V|}} \quad \left( = \mathbb{P}_{x \in [n]^V} [x \in \mathcal{A}] \right).$$

- For  $S \subseteq V$ , the *S-projection density* of  $\mathcal{A}$  is the quantity  $\mu(\text{proj}_S(\mathcal{A}))$ . Note that

$$\mu(\text{proj}_S(\mathcal{A})) = \mathbb{P}_{y \in [n]^S} \left[ \bigvee_{z \in [n]^{V \setminus S}} yz \in \mathcal{A} \right].$$

- The *S-restriction density* of  $\mathcal{A}$ , denoted  $\mu_S(\mathcal{A})$ , is defined by

$$\mu_S(\mathcal{A}) := \max_{z \in [n]^{V \setminus S}} \mu(\mathcal{A}|_S^z) \quad \left( = \max_{z \in [n]^{V \setminus S}} \mathbb{P}_{y \in [n]^V} [yz \in \mathcal{A}] \right).$$

Note that  $\mu_S(\mathcal{A}) \leq \mu(\text{proj}_S(\mathcal{A}))$  and  $\mu_V(\mathcal{A}) = \mu(\mathcal{A})$  and  $\mu_\emptyset(\mathcal{A}) = 1_{[\mathcal{A} \neq \emptyset]}$ .

**Exercise 6.** Let  $\mathcal{A} \subseteq [n]^V$  and  $S \subseteq V$ .

- Show that  $\mu(\mathcal{A}) \leq \mu(\text{proj}_S(\mathcal{A})) \cdot \mu_{V \setminus S}(\mathcal{A})$ .
- Show that inequality (iii) cannot be strengthened to  $\mu(\mathcal{A}) \leq \mu(\text{proj}_S(\mathcal{A})) \cdot \mu(\text{proj}_{V \setminus S}(\mathcal{A}))$ . That is, find an example of  $\mathcal{A}$  and  $S$  where this inequality is false.
- For  $S'' \subseteq S' \subseteq S$ , show that  $\mu_{S''}(\text{proj}_S(\mathcal{A})) \leq \mu_{S''}(\text{proj}_{S'}(\mathcal{A}))$ .

## 3 Bounding the density of a sub-relation of a join

The following lemma will play a role in our lower bound on the  $\text{AC}^0$  formula size of  $\text{SUB}(\text{Path}_k)$ .

**Lemma 7.** Suppose  $\mathcal{A} \subseteq [n]^V$  and  $\mathcal{B} \subseteq [n]^W$  and  $\mathcal{C} \subseteq [n]^{V \cup W}$  such that  $\mathcal{C} \subseteq \mathcal{A} \bowtie \mathcal{B}$ . Then for all  $S \subseteq V$  and  $T \subseteq W$ , we have

$$\mu(\mathcal{C}) \leq \mu(\text{proj}_S(\mathcal{A})) \cdot \mu_{T \setminus S}(\text{proj}_T(\mathcal{B})) \cdot \mu_{(V \cup W) \setminus (S \cup T)}(\mathcal{C}).$$

*Proof.* Lemma 7 is mainly derived by two applications of inequality (iii). First we apply (iii) with respect to  $\mathcal{C}$  and  $S \cup T$ :

$$\mu(\mathcal{C}) \leq \mu(\text{proj}_{S \cup T}(\mathcal{C})) \cdot \mu_{(V \cup W) \setminus (S \cup T)}(\mathcal{C}).$$

Next we apply (iii) with respect to  $\text{proj}_{S \cup T}(\mathcal{C})$  and  $S$ :

$$\mu(\text{proj}_{S \cup T}(\mathcal{C})) \leq \mu(\text{proj}_S(\mathcal{C})) \cdot \mu_{T \setminus S}(\text{proj}_{S \cup T}(\mathcal{C})).$$

Now observe that

$$\begin{aligned} \mu(\text{proj}_S(\mathcal{C})) &\stackrel{(ii)}{\leq} \mu(\text{proj}_S(\mathcal{A})), \\ \mu_{T \setminus S}(\text{proj}_{S \cup T}(\mathcal{C})) &\stackrel{(v)}{\leq} \mu_{T \setminus S}(\text{proj}_T(\mathcal{C})) \stackrel{(ii)}{\leq} \mu_{T \setminus S}(\text{proj}_T(\mathcal{B})). \end{aligned}$$

Combining the four above inequalities finishes the proof.  $\square$