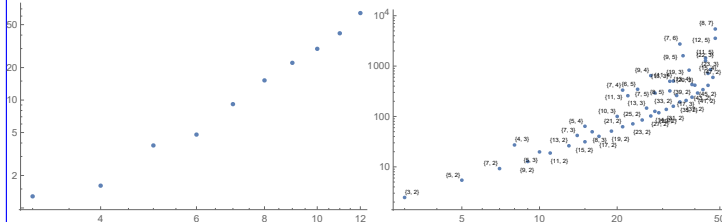




**Abstract.** It has long been known that there are knot invariants associated to semi-simple Lie algebras, and there has long been a dogma as for how to extract them: “quantize and use **representation theory**”. We present an alternative and better procedure: “centrally extend, **approximate by solvable**, and learn how to **re-order exponentials** in a universal enveloping algebra”. While equivalent to the old invariants via a complicated process, our invariants are in practice stronger, faster to compute (poly-time vs. exp-time), and clearly carry topological information.

**KiW 43 Abstract** ([wεβ/kiw](http://wεβ/kiw)). Whether or not you like the formulas on this page, they describe the strongest truly computable knot invariant we know.

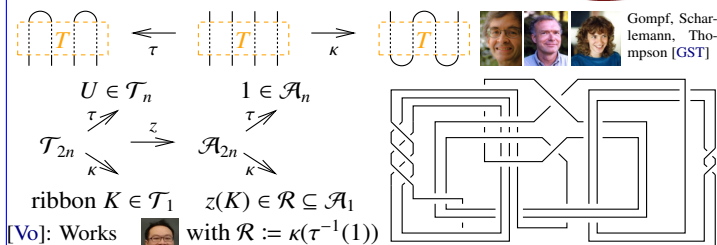
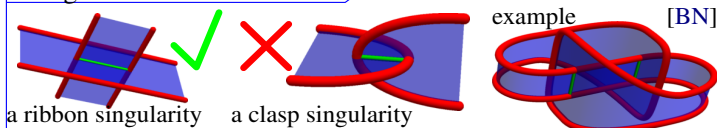
**Experimental Analysis** ([wεβ/Exp](http://wεβ/Exp)). Log-log plots of computation time (sec) vs. crossing number, for all knots with up to 12 crossings (mean times) and for all torus knots with up to 48 crossings:



**Power.** On the 250 knots with at most 10 crossings, the pair  $(\omega, \rho_1)$  attains 250 distinct values, while (Khovanov, HOMFLY-PT) attains only 249 distinct values. To 11 crossings the numbers are (802, 788, 772) and to 12 they are (2978, 2883, 2786).

**Genus.** Up to 12 crossings, always  $\rho_1$  is symmetric under  $t \leftrightarrow t^{-1}$ . With  $\rho_1^+$  denoting the positive-degree part of  $\rho_1$ , always  $\deg \rho_1^+ \leq 2g - 1$ , where  $g$  is the 3-genus of  $K$  (equality for 2530 knots). This gives a lower bound on  $g$  in terms of  $\rho_1$  (conjectural, but undoubtedly true). This bound is often weaker than the Alexander bound, yet for 10 of the 12-xing Alexander failures it does give the right answer.

**Ribbon Knots.**



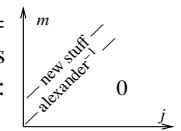
[Vo]: Works with  $\mathcal{R} := \kappa(\tau^{-1}(1))$  for Alexander!  
 $A^+ = -t^8 + 2t^7 - t^6 - 2t^4 + 5t^3 - 2t^2 - 7t + 13$   
 $\rho_1^+ = 5t^{15} - 18t^{14} + 33t^{13} - 32t^{12} + 2t^{11} + 42t^{10} - 62t^9 - 8t^8 + 166t^7 - 242t^6 +$   
 Faster is better, leaner is meaner!  $108t^5 + 132t^4 - 226t^3 + 148t^2 - 11t - 36$

**Ordering Symbols.**  $\odot$  (*poly* | *specs*) plants the variables of *poly* in  $S(\otimes_i \mathfrak{g})$  on several tensor copies of  $\mathcal{U}(\mathfrak{g})$  according to *specs*. E.g.,  
 $\odot(a_1^3 y_1 a_2 e^{y_3} x_3^9 | x_3 a_1 \otimes y_1 y_3 a_2) = x^9 a^3 \otimes y e^y a \in \mathcal{U}(\mathfrak{g}) \otimes \mathcal{U}(\mathfrak{g})$   
 This enables the description of elements of  $\hat{\mathcal{U}}(\mathfrak{g})^{\otimes S}$  using commutative polynomials / power series.

**Theorem** ([BNG], conjectured [MM], elucidated [Ro1]). Let  $J_d(K)$  be the coloured Jones polynomial of  $K$ , in the  $d$ -dimensional representation of  $sl_2$ . Writing

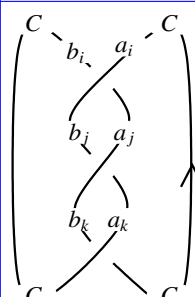
$$\left. \frac{(q^{1/2} - q^{-1/2})J_d(K)}{q^{d/2} - q^{-d/2}} \right|_{q=e^h} = \sum_{j,m \geq 0} a_{jm}(K) d^j h^m,$$

“below diagonal” coefficients vanish,  $a_{jm}(K) = 0$  if  $j > m$ , and “on diagonal” coefficients give the inverse of the Alexander polynomial:  
 $(\sum_{m=0}^{\infty} a_{mm}(K) h^m) \cdot \omega(K)(e^h) = 1.$



“Above diagonal” we have **Rozansky’s Theorem** [Ro3, (1.2)]:

$$J_d(K)(q) = \frac{q^d - q^{-d}}{(q - q^{-1})\omega(K)(q^d)} \left( 1 + \sum_{k=1}^{\infty} \frac{(q-1)^k \rho_k(K)(q^d)}{\omega^{2k}(K)(q^d)} \right).$$



**The Yang-Baxter Technique.** Given an algebra  $U$  (typically  $\hat{\mathcal{U}}(\mathfrak{g})$  or  $\hat{\mathcal{U}}_q(\mathfrak{g})$ ) and elements

$$R = \sum a_i \otimes b_i \in U \otimes U \quad \text{and} \quad C \in U,$$

form

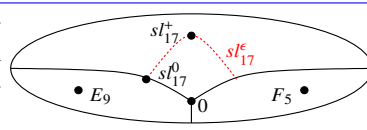
$$Z = \sum_{i,j,k} C a_i b_j a_k C^2 b_i a_j b_k C.$$

**Problem.** Extract information from  $Z$ .

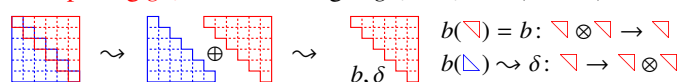
**The Dogma.** Use representation theory. In principle finite, but *slow*.

**The Loyal Opposition.** For certain algebras, work in a homomorphic poly-dimensional “space of formulas”.  
 $m_k^{ij} \circlearrowleft \{ \mathcal{F}_S \} \xrightarrow{\mathbb{E}} \{ U^{\otimes S} \} \circlearrowright m_k^{ij}$

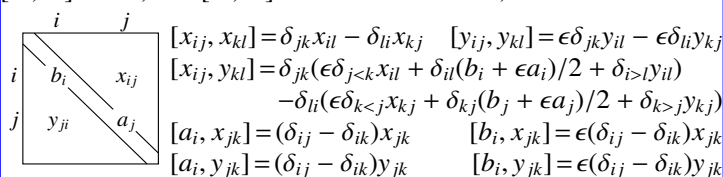
**The (fake) moduli** of Lie algebras on  $V$ , a quadratic variety in  $(V^*)^{\otimes 2} \otimes V$  is on the right. We care about  $sl_{17}^k := sl_{17}^e / (\epsilon^{k+1} = 0)$ .



**Recomposing  $gl_n$ .** Half is enough!  $gl_n \oplus \mathfrak{a}_n = \mathcal{D}(\nabla, b, \delta)$ :



Now define  $gl_n^e := \mathcal{D}(\nabla, b, \epsilon\delta)$ . Schematically, this is  $[\nabla, \nabla] = \nabla$ ,  $[\Delta, \Delta] = \epsilon\Delta$ , and  $[\nabla, \Delta] = \Delta + \epsilon\nabla$ . In detail, it is



**The Main  $sl_2$  Theorem.** Let  $\mathfrak{g}^e = \langle t, y, a, x \rangle / ([t, \cdot] = 0, [a, x] = x, [a, y] = -y, [x, y] = t - 2\epsilon a)$  and let  $\mathfrak{g}_k = \mathfrak{g}^e / (\epsilon^{k+1} = 0)$ . The  $\mathfrak{g}_k$ -invariant of any  $S$ -component tangle  $K$  can be written in the form  $Z(K) = \odot(\omega e^{L+Q+P} : \otimes_{i \in S} y_i a_i x_i)$ , where  $\omega$  is a scalar (a rational function in the variables  $t_i$  and their exponentials  $T_i := e^{t_i}$ ), where  $L = \sum l_{ij} t_i a_j$  is a quadratic in  $t_i$  and  $a_j$  with integer coefficients  $l_{ij}$ , where  $Q = \sum q_{ij} y_i x_j$  is a quadratic in the variables  $y_i$  and  $x_j$  with scalar coefficients  $q_{ij}$ , and where  $P$  is a polynomial in  $\{\epsilon, y_i, a_i, x_i\}$  (with scalar coefficients) whose  $\epsilon^d$ -term is of degree at most  $2d + 2$  in  $\{y_i, \sqrt{a_i}, x_i\}$ . Furthermore, after setting  $t_i = t$  and  $T_i = T$  for all  $i$ , the invariant  $Z(K)$  is poly-time computable.