## 2 - Integers

Jean-Baptiste Campesato

In this chapter, we are going to construct the set $\mathbb{Z}$ of integers and then to study its properties. The informal idea consists in extending $\mathbb{N}$ by adding its symmetry with respect to 0 :


For this purpose, we have to give a meaning to the notation $-n$ where $n$ is a natural number and then we have to extend from $\mathbb{N}$ to $\mathbb{Z}$ the operations $(+, \times)$ and the order $(\leq)$.

There are several ways to formally do that. The usual one consists in defining $\mathbb{Z}=(\mathbb{N} \times \mathbb{N}) / \sim$ for the equivalence relation $(a, b) \sim(c, d) \Leftrightarrow a+d=b+c$. Let me explain what does it mean: intuitively $(a, b)$ stands for $a-b$, but, since such an expression is not unique (e.g. $7-5=10-8$ ), we need to "identify" some couples (e.g. $(7,5)=(10,8)$ ). This construction has several advantages (it is easy to extend,$+ \times$ and $\leq$ ) but it needs an additional layer of abstraction (equivalence relations, equivalence classes...).

Instead, I will use a more naive approach. The counterpart is that extending the operations will be a little bit tedious with several cases to handle (e.g. the definition of $a+b$ will depend on the signs of $a$ and $b$, so we have 4 cases just to define the addition...).

Note that what we are going to describe in a few lines took centuries to be developped and accepted: during the 18th century, most mathematicians were still reluctant about using negative numbers.

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## 1 Construction of the integers

### 1.1 Definition

Definition 1. For any $n \in \mathbb{N} \backslash\{0\}$, we formally introduce the symbol $-n$ read as minus $n$ and we fix the convention that $-0=0$.
We define the set $-\mathbb{N}:=\{-n: n \in \mathbb{N}\}$. Then the set of integers is $\mathbb{Z}:=(-\mathbb{N}) \cup \mathbb{N}$.
Remark 2. $(-\mathbb{N}) \cap \mathbb{N}=\{0\}$
Remark 3. $\mathbb{N} \subset \mathbb{Z}$


### 1.2 Operations

Definition 4. For $m, n \in \mathbb{N}$, we set:
(i) $m+n$ for the usual addition in $\mathbb{N}$
(ii) $(-m)+(-n)=-(m+n)$
(iii) $m+(-n)=\left\{\begin{array}{c}k \text { where } k \text { is the unique natural integer such that } m=n+k \text { if } n \leq m \\ -k \text { where } k \text { is the unique natural integer such that } n=m+k \text { if } m \leq n\end{array}\right.$
(iv) $(-m)+n=n+(-m)$ where $n+(-m)$ is defined in (iii)

We've just defined $+: \begin{array}{ccc}\mathbb{Z} \times \mathbb{Z} & \rightarrow & \mathbb{Z} \\ (a, b) & \mapsto & a+b\end{array}$.
Remark 5. We have to check that the overlapping cases $m=0$ or $n=0$ are not contradictory.
Definition 6. For $m, n \in \mathbb{N}$, we set:
(i) $m \times n$ for the usual product in $\mathbb{N}$
(ii) $(-m) \times(-n)=m \times n$
(iii) $m \times(-n)=-(m \times n)$
(iv) $(-m) \times n=-(m \times n)$

We've just defined $\times: \begin{array}{ccc}\mathbb{Z} \times \mathbb{Z} & \rightarrow & \mathbb{Z} \\ (a, b) & \mapsto & a \times b\end{array}$.
Remark 7. We may simply write $a b$ for $a \times b$ when there is no possible confusion.
Remark 8. Note that the addition and product on $\mathbb{Z}$ are compatible with the addition and product on $\mathbb{N}$.
Definition 9. For $n \in \mathbb{N}$, we set $-(-n)=n$. Then $-a$ is well-defined for every $a \in \mathbb{Z}$.

## Proposition 10.

$\bullet+$ is associative: $\forall a, b, c \in \mathbb{Z},(a+b)+c=a+(b+c)$

- 0 is the unit of $+: \forall a \in \mathbb{Z}, a+0=0+a=a$
- $-a$ is the additive inverse of $a: \forall a \in \mathbb{Z}, a+(-a)=(-a)+a=0$
$\bullet$ + is commutative: $\forall a, b \in \mathbb{Z}, a+b=b+a$
- $\times$ is associative: $\forall a, b, c \in \mathbb{Z},(a b) c=a(b c)$
- $\times$ is distributive with respect to $+: \forall a, b, c \in \mathbb{Z}, a \times(b+c)=a b+a c$ et $(a+b) c=a c+b c$
- 1 is the unit of $\times: \forall a \in \mathbb{Z}, 1 \times a=a \times 1=a$
- $\times$ is commutative: $\forall a, b \in \mathbb{Z}, a b=b a$
- $\forall a, b \in \mathbb{Z}, a b=0 \Rightarrow(a=0$ or $b=0)$

The above properties are easy to prove but the proofs are tedious with several cases depending on the signs.

Remark. From now on, we may simply write $a-b$ for $a+(-b)$ and $-a+b$ for $(-a)+b$.
Corollary 11. $\forall a, b, c \in \mathbb{Z},(a c=b c$ and $c \neq 0) \Longrightarrow a=b$
Proof. Let $a, b, c \in \mathbb{Z}$ be such that $a c=b c$ and $c \neq 0$.
Then $(a-b) c=0$. So either $a-b=0$ or $c=0$. Since $c \neq 0$, we get $a-b=0$, i.e. $a=b$..

### 1.3 Order

Definition 12. We define the binary relation $\leq$ on $\mathbb{Z}$ by

$$
\forall a, b \in \mathbb{Z}, a \leq b \Leftrightarrow b-a \in \mathbb{N}
$$

Proposition 13. $\leq$ defines a total order on $\mathbb{Z}$.
Proof.

- Reflexivity. Let $a \in \mathbb{Z}$, then $a-a=0 \in \mathbb{N}$ so $a \leq a$.
- Antisymmetry. Let $a, b \in \mathbb{Z}$. Assume that $a \leq b$ and that $b \leq a$. Then $b-a \in \mathbb{N}$ and $a-b \in \mathbb{N}$. So $a-b=-(b-a) \in(-\mathbb{N})$. Hence $a-b \in(-\mathbb{N}) \cap \mathbb{N}=\{0\}$ and thus $a=b$.
- Transitivity. Let $a, b, c \in \mathbb{Z}$. Assume that $a \leq b$ and that $b \leq c$. Then $b-a \in \mathbb{N}$ and $c-b \in \mathbb{N}$.

Thus $c-a=(c-b)+(b-a) \in \mathbb{N}$, i.e. $a \leq c$.

- Let $a, b \in \mathbb{Z}$. Then $b-a \in \mathbb{Z}=(-\mathbb{N}) \cup(\mathbb{N})$.

First case: $b-a \in \mathbb{N}$ then $a \leq b$.
Second case: $b-a \in(-\mathbb{N})$, then $a-b=-(b-a) \in \mathbb{N}$ and $b \leq a$.
Hence the order is total.

Proposition 14. The order on $\mathbb{Z}$ is compatible with the order on $\mathbb{N}$.
Proof. Let $a, b \in \mathbb{N}$.

- Assume that $a \leq_{\mathbb{Z}} b$. Then $k=b-a \in \mathbb{N}$. So $b=a+k$, i.e. $a \leq_{\mathbb{N}} b$.
- Assume that $a \leq_{\mathbb{N}} b$. Then $b=a+k$ for some $k \in \mathbb{N}$. Then $b-a=k \in \mathbb{N}$, i.e. $a \leq_{\mathbb{Z}} b$.


## Proposition 15.

1. $\mathbb{N}=\{a \in \mathbb{Z}, 0 \leq a\}$
2. $\forall a, b, c \in \mathbb{Z}, a \leq b \Leftrightarrow a+c \leq b+c$
3. $\forall a, b, c, d \in \mathbb{Z},(a \leq b$ and $c \leq d) \Rightarrow a+c \leq b+d$
4. $\forall a, b \in \mathbb{Z}, \forall c \in \mathbb{N} \backslash\{0\}, a \leq b \Leftrightarrow a c \leq b c$
5. $\forall a, b \in \mathbb{Z}, \forall c \in(-\mathbb{N}) \backslash\{0\}, a \leq b \Leftrightarrow b c \leq a c$

Proof.

1. Let $a \in \mathbb{Z}$. Then $0 \leq a \Leftrightarrow a=a-0 \in \mathbb{N}$.
2. Let $a, b, c \in \mathbb{Z}$. Then $a \leq b \Leftrightarrow b-a \in \mathbb{N} \Leftrightarrow(b+c)-(a+c) \in \mathbb{N} \Leftrightarrow a+c \leq b+c$.
3. Let $a, b, c, d \in \mathbb{Z}$. Assume that $a \leq b$ and that $c \leq d$. Then $b-a \in \mathbb{N}$ and $d-c \in \mathbb{N}$. Hence $(b+d)-(a+c)=(b-a)+(d-c) \in \mathbb{N}$, i.e. $a+c \leq b+d$.
4. Let $a, b \in \mathbb{Z}$ and $c \in \mathbb{N}$.
$\Rightarrow$ : Assume that $a \leq b$. Then $b-a \in \mathbb{N}$, thus $b c-a c=(b-a) c \in \mathbb{N}$. Therefore $a c \leq b c$.
$\Leftarrow$ : Assume that $c \neq 0$ and that $a c \leq b c$. Then $b c-a c=(b-a) c \in \mathbb{N}$. Assume by contradiction that $(b-a) \in(-\mathbb{N}) \backslash\{0\}$ then, by definition of the multiplication, $(b-a) c \in(-\mathbb{N}) \backslash\{0\}$, which is a contradiction. Hence $b-a \in \mathbb{N}$, i.e. $a \leq b$.
5. Let $a, b \in \mathbb{Z}$ and $c \in(-\mathbb{N})$.
$\Rightarrow$ : Assume that $a \leq b$. Then $b-a \in \mathbb{N}$, thus $a c-b c=(b-a)(-c) \in \mathbb{N}$. Therefore $b c \leq a c$.
$\Leftarrow$ : Assume that $c \neq 0$ and that $b c \leq a c$. Then $a c-b c=(b-a)(-c) \in \mathbb{N}$. And we conclude as in 4 .

Remark 16. Given $a, b, c \in \mathbb{Z}$, it is common to lighten the notation by writing $a \leq b \leq c$ for $(a \leq b$ and $b \leq c)$.

## Theorem 17.

1. A non-empty subset $A$ of $\mathbb{Z}$ which is bounded from below has a least element, i.e.

$$
\exists m \in A, \forall a \in A, m \leq a
$$

2. A non-empty subset $A$ of $\mathbb{Z}$ which is bounded from above has a greatest element, i.e.

$$
\exists M \in A, \forall a \in A, a \leq M
$$

Proof.

1. Assume that $A$ is a non-empty subset of $\mathbb{Z}$ which is bounded from below.

Then there exists $k \in \mathbb{Z}$ such that $\forall a \in A, k \leq a$. Define $S=\{a-k: a \in A\}$.
Then $S$ is a non-empty subset of $\mathbb{N}$ (indeed, $\forall a \in A, 0 \leq a-k)$.
By the well-ordering principle, there exists $\tilde{m} \in S$ such that $\forall a \in A, \tilde{m} \leq a-k$.
Then $m=\tilde{m}+k$ is the least element of $A$ (note that $\tilde{m} \in S$ so $m=\tilde{m}+k \in A$ )
2. Assume that $A$ is a non-empty subset of $\mathbb{Z}$ which is bounded from above. Then $(-A)=\{-a: a \in A\}$ is a non-empty subset of $\mathbb{Z}$ which is bounded from below (prove it).
By the above, there exists $m \in(-A)$ such that $\forall a \in A, m \leq-a$. Hence $\forall a \in A, a \leq-m$.
Thus $M:=-m$ is the greatest element of $A$.

## 2 Absolute value

Definition 18. For $n \in \mathbb{Z}$, we define the absolute value of $n$ by $|n|:=\left\{\begin{array}{cl}n & \text { if } n \in \mathbb{N} \\ -n & \text { if } n \in(-\mathbb{N})\end{array}\right.$.
Proposition 19.
(i) $\forall n \in \mathbb{Z},|n| \in \mathbb{N}$
(ii) $\forall n \in \mathbb{Z}, n \leq|n|$
(iii) $\forall n \in \mathbb{Z},|n|=0 \Leftrightarrow n=0$
(iv) $\forall a, b \in \mathbb{Z},|a b|=|a||b|$
(v) $\forall a, b \in \mathbb{Z},|a| \leq b \Leftrightarrow-b \leq a \leq b$
(vi) $\forall a, b \in \mathbb{Z},|a+b| \leq|a|+|b| \quad$ (triangle inequality)

Proof.
(i) First case: if $n \in \mathbb{N}$ then $|n|=n \in \mathbb{N}$.

Second case: if $n \in(-\mathbb{N})$ then $n=-m$ for some $m \in \mathbb{N}$ and $|n|=-n=-(-m)=m \in \mathbb{N}$.
(ii) First case: $n \in \mathbb{N}$. Then $n \leq n=|n|$.

Second case: $n \in(-\mathbb{N})$. Then $n \leq 0 \leq|n|$.
(iii) Note that $|0|=0$ and that if $n \neq 0$ then $|n| \neq 0$.
(iv) You have to study separately the four cases depending on the signs of $a$ and $b$.
(v) If $b<0$ then $|a| \leq b$ and $-b \leq a \leq b$ are both false. So we may assume that $b \in \mathbb{N}$. Then

First case: $a \in \mathbb{N}$. Then $|a| \leq b \Leftrightarrow a \leq b \Leftrightarrow-b \leq a \leq b$.
Second case: $a \in(-\mathbb{N})$. Then $|a| \leq b \Leftrightarrow-a \leq b \Leftrightarrow-b \leq a \Leftrightarrow-b \leq a \leq b$.
(vi) Since $a+b \leq|a|+|b|$ and $-(a+b)=-a-b \leq|-a|+|-b|=|a|+|b|$, we get $|a+b| \leq|a|+|b|$.

## 3 Euclidean division

Theorem 20 (Euclidean division).
Given $a \in \mathbb{Z}$ and $b \in \mathbb{Z} \backslash\{0\}$, there exists a unique couple $(q, r) \in \mathbb{Z}^{2}$ such that

$$
\left\{\begin{array}{l}
a=b q+r \\
0 \leq r<|b|
\end{array}\right.
$$

The integers $q$ and $r$ are respectively called the quotient and the remainder of the division of $a$ by $b$.

Proof.

## Existence:

First case: assume that $0<b$.
We set ${ }^{1} E=\{p \in \mathbb{Z}: b p \leq a\}$.

- $E \neq \varnothing$, indeed if $0 \leq a$ then $0 \in E$, otherwise $a \in E$.
- $|a|$ is an upper bound of $E$.

Indeed, let $p \in E$.
If $p \leq 0$ then $p \leq 0 \leq|a|$.
Otherwise, if $0<p$ then $1 \leq b \Longrightarrow p \leq b p \leq a \leq|a|$.
Thus $E$ is a non-empty subset of $\mathbb{Z}$ which is bounded from above.
Hence it admits a greatest element, i.e. there exists $q \in E$ such that $\forall p \in E, p \leq q$.
We set $r=a-b q$. Since $q \in E, r=a-b q \geq 0$.
And $q+1 \notin E$ since $q+1>q$ whereas $q$ is the greatest element of $E$.
Therefore $b(q+1)>a$, so $r=a-b q<b=|b|$.
We wrote $a=b q+r$ with $0 \leq r<|b|$ as expected.
Second case: assume that $b<0$.
Then we apply the first case to $a$ and $-b>0$ : there exists $(q, r) \in \mathbb{Z}^{2}$ such that $a=-b q+r=b(-q)+r$ with $0 \leq r<-b=|b|$.

Uniqueness: assume that we have two suitable couples ( $q, r$ ) and $\left(q^{\prime}, r^{\prime}\right)$.
Then $r^{\prime}-r=\left(a-b q^{\prime}\right)-(a-b q)=b\left(q-q^{\prime}\right)$. Besides

$$
\left\{\begin{array} { l } 
{ 0 \leq r < | b | } \\
{ 0 \leq r ^ { \prime } < | b | }
\end{array} \Longrightarrow \left\{\begin{array}{l}
-|b|<-r \leq 0 \\
0 \leq r^{\prime}<|b|
\end{array} \Longrightarrow-|b|<r^{\prime}-r<|b|\right.\right.
$$

Thus $-|b|<b\left(q-q^{\prime}\right)<|b|$, from which we get $|b|\left|q-q^{\prime}\right|=\left|b\left(q-q^{\prime}\right)\right|<|b|$.
Since $|b|>0$, we obtain $0 \leq\left|q-q^{\prime}\right|<1$.
But we proved in the first chapter that there is no natural number between 0 and 1 .
Therefore $\left|q-q^{\prime}\right|=0$, which implies that $q-q^{\prime}=0$, i.e. $q=q^{\prime}$.
Finally, $r^{\prime}=b-a q^{\prime}=b-a q=r$.

## Examples 21.

- Division of 22 by 5: $22=5 \times 4+2$.

The quotient is $q=4$ and the remainder is $r=2$.

- Division of -22 by $5:-22=5 \times(-5)+3$.

The quotient is $q=-5$ and the remainder is $r=3$.

- Division of 22 by -5 : $22=(-5) \times(-4)+2$.

The quotient is $q=-4$ and the remainder is $r=2$.

- Division of -22 by -5 : $-22=(-5) \times 5+3$.

The quotient is $q=5$ and the remainder is $r=3$.
Proposition 22. Given $n \in \mathbb{Z}$,

- either $n=2 k$ for some $k \in \mathbb{Z}$ (then we say that $n$ is even),
- or $n=2 k+1$ for some $k \in \mathbb{Z}$ (then we say that $n$ is odd),
and these cases are exclusive.
Proof. Let $n \in \mathbb{Z}$. By the Euclidean division by 2 , there exist $k, r \in \mathbb{Z}$ such that $n=2 k+r$ and $0 \leq r \leq 1$. But we know from the last chapter that there is no natural number between 0 and 1 . Hence either $r=0$ or $r=1$. These cases are exclusive by the uniqueness of the Euclidean division.

[^0]
## 4 Divisibility

Definition 23. Given $a, b \in \mathbb{Z}$, we say that $a$ is divisible by $b$ if there exists $k \in \mathbb{Z}$ such that $a=b k$. In this case we write $b \mid a$ and we also say that $b$ is divisor of $a$ or that $a$ is a multiple of $b$.

Examples 24. • $(-5) \mid 10 \bullet 5 \nmid(-11)$
(we will study divisibility criteria later in the term)

## Remarks 25.

- Any integer is a divisor of 0 , i.e $\forall b \in \mathbb{Z}, b \mid 0$. Indeed, $0=b \times 0$.
- Any integer is divisible by 1 and itself, i.e. $\forall a \in \mathbb{Z}, 1 \mid a$ and $a \mid a$. Indeed, $a=1 \times a=a \times 1$.
- The only integer divisible by 0 is 0 itself, i.e. $\forall a \in \mathbb{Z}, 0 \mid a \Longrightarrow a=0$. Indeed, then $a=0 \times k$ for some $k \in \mathbb{Z}$ and hence $a=0$.
- When $b \neq 0, b \mid a$ if and only if the remainder of the Euclidean division of $a$ by $b$ is $r=0$.


## Proposition 26.

1. $\forall a, b \in \mathbb{Z},(a \mid b$ and $b \mid a) \Longrightarrow|a|=|b|$
2. $\forall a, b, c \in \mathbb{Z},(a \mid b$ and $b \mid c) \Longrightarrow a \mid c$
3. $\forall a, b, c, d \in \mathbb{Z},(a \mid b$ and $c \mid d) \Longrightarrow a c \mid b d$
4. $\forall a, b, c, \lambda, \mu \in \mathbb{Z},(a \mid b$ and $a \mid c) \Longrightarrow a \mid(\lambda b+\mu c)$
5. $\forall a \in \mathbb{Z}, a|1 \Longrightarrow| a \mid=1$

Proof.

1. Let $a, b \in \mathbb{Z}$ satisfying $a \mid b$ and $b \mid a$. If $a=0$ then $b=0$ (from $0 \mid b$ ). So we may assume that $a \neq 0$.

There exist $k, l \in \mathbb{Z}$ such that $b=a k$ and $a=b l$. Then $a=b l=a k l$, thus $1=k l$ since $a \neq 0$.
Therefore, $1=|1|=|k l|=|k| \times|l|$. Since $|k|,|l| \in \mathbb{N}$, we get that $|k|=|l|=1$.
Finally, $|a|=|b l|=|b| \times|l|=|b| \times 1=|b|$.
2. Let $a, b, c \in \mathbb{Z}$ satisfying $a \mid b$ and $b \mid c$. Then $b=a k$ and $c=b l$ for some $k, l \in \mathbb{Z}$.

Therefore $c=b l=a k l$, so $a \mid c$.
3. Let $a, b, c, d \in \mathbb{Z}$ satisfying $a \mid b$ and $c \mid d$. Then $b=a k$ and $d=c l$ for some $k, l \in \mathbb{Z}$.

Therefore $b d=a c k l$, so $a c \mid b d$.
4. Let $a, b, c \in \mathbb{Z}$ satisfying $a \mid b$ and $a \mid c$. Then $b=k a$ and $c=l a$ for some $k, l \in \mathbb{Z}$.

Hence $\lambda b+\mu c=\lambda k a+\mu l a=(\lambda k+\mu l) a$. Thus $a \mid(\lambda b+\mu c)$.
5. Let $a \in \mathbb{Z}$. Assume that $a \mid 1$. Then $a \mid 1$ and $1 \mid a$. So by the first item, $|a|=1$.

## 5 Greatest common divisor

Theorem 27. Given $a, b \in \mathbb{Z}$ not both zero, the set common divisors of $a$ and $b$ admits a greatest element denoted $\operatorname{gcd}(a, b)$ and called the greatest common divisor of $a$ and $b$.

Proof. Let $a, b \in \mathbb{Z}$ not both zero. We set $S=\{d \in \mathbb{Z}: d \mid a$ and $d \mid b\}$.

- $S$ is non-empty since it contains 1.
- Since $a$ and $b$ are not both zero, we know that $a \neq 0$ or $b \neq 0$.

Without loss of generality, let assume that $a \neq 0$.
Let $d \in S$ then $a=d k$ for some $k \in \mathbb{Z}$. Note that $k \neq 0$ (otherwise $a=d k=0$ ), hence $1 \leq|k|$.
Thus $d \leq|d| \leq|d| \times|k|=|d k|=|a|$. Hence $S$ is bounded from above by $|a|$.
Therefore, $S$ admits a greatest element (as an non-empty subset of $\mathbb{Z}$ bounded from above).
Remark 28. Note that $\operatorname{gcd}(a, b) \geq 1$ since 1 is a common divisor of $a$ and $b$ (particularly $\operatorname{gcd}(a, b) \in \mathbb{N})$.
Proposition 29. Let $a, b \in \mathbb{Z}$ not both zero and $d \in \mathbb{N} \backslash\{0\}$. Then

$$
\left\{\left.\begin{array}{l}
d \mid a \\
d \mid b \\
\forall \delta \in \mathbb{N},(\delta \mid \text { a and } \delta \mid b) \Longrightarrow \delta \mid d
\end{array} \right\rvert\, \Longrightarrow d=\operatorname{gcd}(a, b)\right.
$$

Remark 30. We will see later that the converse holds (Proposition 35.(3)).
Proof. Let $a, b \in \mathbb{Z}$ not both zero and $d \in \mathbb{N} \backslash\{0\}$. Assume that $d|a, d| b$ and that $d$ is a multiple of every non-negative common divisors, i.e.

$$
\forall \delta \in \mathbb{N},(\delta \mid a \text { and } \delta \mid b) \Longrightarrow \delta \mid d
$$

Then $d$ is a common divisor of $a$ and $b$. We need to prove that it is the greatest one.
Let $\delta \in \mathbb{Z}$ be a common divisor of $a$ and $b$.

- If $\delta \leq 0$ then $\delta \leq d$.
- If $\delta>0$ then $d=\delta k$ for some $k \in \mathbb{Z}$.

Note that $k \geq 1$ since $d, \delta>0$. Thus $\delta \leq \delta k=d$

The following theorem is extremely useful! We will use it quite often to study gcd and also when studying modular arithmetic!

Theorem 31 (Bézout's identity). Given $a, b \in \mathbb{Z}$ not both zero, there exist $u, v \in \mathbb{Z}$ such that

$$
a u+b v=\operatorname{gcd}(a, b)
$$

Example 32. $\operatorname{gcd}(15,25)=5=15 \times 2+25 \times(-1)$.
We will see below an algorithm in order to find a suitable couple $(u, v)$.

## Remarks 33.

- The couple $(u, v)$ is not unique: $5=15 \times 27+25 \times(-16)$.
- The converse is false: $2=3 \times 4+5 \times(-2)$ but $\operatorname{gcd}(3,5)=1 \neq 2$.

Nonetheless, we will see later that there is a partial converse when $\operatorname{gcd}(a, b)=1$.
Proof of Theorem 31. Let $a, b \in \mathbb{Z}$ not both zero. Set $S=\{n \in \mathbb{N} \backslash\{0\}: \exists u, v \in \mathbb{Z}, n=a u+b v\}$.
Without loss of generality we may assume that $a \neq 0$.
Note that $S$ is not empty. Indeed,

- If $a<0$ then $n=a \times(-1)+b \times 0$ is in $S$, or,
- If $a>0$ then $n=a \times 1+b \times 0$ is in $S$.

Thus, by the well-ordering principle, $S$ admits a least element $d$.
Since $d \in S$, we know that $d=a u+b v$ for some $u, v \in \mathbb{Z}$.
Let's prove that $d=\operatorname{gcd}(a, b)$.

- By Euclidean division, there exist $q, r \in \mathbb{Z}$ such that $a=d q+r$ and $0 \leq r<|d|=d$.

Assume by contradiction that $r \neq 0$.
Then $r=a-q d=a-q(a u+b v)=a \times(1-q u)+b \times(-q v)$ is in $S$. Which contradicts the fact that $d$ is the least element of $S$. Hence $r=0$ and $a=d q$, i.e. $d \mid a$.

- Similarly $d \mid b$.
- Let $\delta \in \mathbb{Z}$ be another common divisor of $a$ and $b$.

Since $\delta \mid a$ and $\delta \mid b, a=\delta k$ and $b=\delta l$ for some $k, l \in \mathbb{Z}$. Hence $d=a u+b v=\delta(k u+l v)$, i.e. $\delta \mid d$.
Therefore, by Proposition 29, $d=\operatorname{gcd}(a, b)$.
Hence $\operatorname{gcd}(a, b)=a u+b v$ as requested.
Proposition 34. $\forall a \in \mathbb{Z} \backslash\{0\}, \operatorname{gcd}(a, 0)=|a|$
Proof. By definition, $\operatorname{gcd}(a, 0)$ is the greatest divisor of $a$.
Since $a=|a| \times( \pm 1)$, we know that $|a|$ is a divisor of $a$. We have to check that it is the greatest one.
Let $d$ be a non-negative divisor of $a$, then $a=d k$ for some $k \in \mathbb{Z}$.
Since $a \neq 0$, we know that $k \neq 0$.
Hence $1 \leq|k|$ from which we get that $d \leq d|k|=|d| \times|k|=|d k|=|a|$.

Proposition 35. Let $a, b \in \mathbb{Z}$ not both zero, then

1. $\operatorname{gcd}(a, b)=\operatorname{gcd}(b, a)$
2. $\operatorname{gcd}(a, b)=\operatorname{gcd}(a,-b)=\operatorname{gcd}(-a, b)=\operatorname{gcd}(-a,-b)$
3. $\forall \delta \in \mathbb{Z},(\delta \mid a$ and $\delta \mid b) \Longrightarrow \delta \mid \operatorname{gcd}(a, b)$
4. $\forall \lambda \in \mathbb{Z} \backslash\{0\}, \operatorname{gcd}(\lambda a, \lambda b)=|\lambda| \operatorname{gcd}(a, b)$
5. $\forall k \in \mathbb{Z}, \operatorname{gcd}(a+k b, b)=\operatorname{gcd}(a, b)$

Proof. I will just prove 3, 4 and 5, the first two being easy to prove.
3. Let $a, b \in \mathbb{Z}$. Let $\delta \in \mathbb{Z}$. Assume that $\delta \mid a$ and $\delta \mid b$.

By Bézout's theorem, $\operatorname{gcd}(a, b)=a u+b v$ for some $u, v \in \mathbb{Z}$.
Since $\delta \mid a$ and $\delta \mid b$, we have that $\delta \mid a u+b v=\operatorname{gcd}(a, b)$.
4. Let $a, b \in \mathbb{Z}$ let $\lambda \in \mathbb{Z} \backslash\{0\}$. Since $|\lambda|$ divides $\lambda a$ and $\lambda b$, then it divides $\operatorname{gcd}(\lambda a, \lambda b)$ by the third item. Hence $\operatorname{gcd}(\lambda a, \lambda b)=|\lambda| \times d$ for some $d \in \mathbb{Z}$. Let's prove that $d=\operatorname{gcd}(a, b)$.
Let $n \in \mathbb{Z}$, then $n|a, b \Leftrightarrow| \lambda|n| \lambda a, \lambda b \Leftrightarrow|\lambda| n|\operatorname{gcd}(\lambda a, \lambda b) \Leftrightarrow n| d$.
5. Let $a, b, k \in \mathbb{Z} . \operatorname{gcd}(a, b) \mid a, b$ hence $\operatorname{gcd}(a, b) \mid a+k b$. Thus $\operatorname{gcd}(a, b) \mid \operatorname{gcd}(a+k b, b)$.

Similarly, $\operatorname{gcd}(a+k b, b) \mid a+k b, b$ hence $\operatorname{gcd}(a+k b, b) \mid a+k b-k b=a$. Thus $\operatorname{gcd}(a+k b, b) \mid \operatorname{gcd}(a, b)$. Hence $|\operatorname{gcd}(a+k b, b)|=|\operatorname{gcd}(a, b)|$. Since they are both non-negative, we get $\operatorname{gcd}(a+k b, b)=\operatorname{gcd}(a, b)$.

## 6 Euclid's algorithm

Euclid's algorithm is an efficient way to compute the gcd of two numbers.

Let $a, b \in \mathbb{Z}$ not both zero.
Initialization of the algorithm. We set $a_{0}:=|a|$ and $b_{0}:=|b|$. Note that $\operatorname{gcd}\left(a_{0}, b_{0}\right)=\operatorname{gcd}( \pm a, \pm b)=\operatorname{gcd}(a, b)$. Iteration. Assume that $a_{n}, b_{n} \in \mathbb{Z}$ are already constructed with $a_{n}, b_{n} \geq 0$ not both zero.

- If $b_{n}=0$ then $\operatorname{gcd}\left(a_{n}, b_{n}\right)=a_{n}$ and the algorithm stops.
- Otherwise, by Euclidean division, there exist $q_{n}, r_{n} \in \mathbb{R}$ such that $a_{n}=b_{n} q_{n}+r_{n}$ and $0 \leq r_{n}<b_{n}$.

We set $a_{n+1}:=b_{n}$ and $b_{n+1}:=r_{n}$, then $a_{n+1}=b_{n}>0$ and $0 \leq b_{n+1}<b_{n}$.
Moreover, using Proposition 35.(5),

$$
\operatorname{gcd}\left(a_{n}, b_{n}\right)=\operatorname{gcd}\left(b_{n} q_{n}+r_{n}, b_{n}\right)=\operatorname{gcd}\left(r_{n}, b_{n}\right)=\operatorname{gcd}\left(b_{n}, r_{n}\right)=\operatorname{gcd}\left(a_{n+1}, b_{n+1}\right)
$$

We repeat the iterative process with $a_{n+1}$ and $b_{n+1}$.
Conclusion. Since the $b_{n}$ are natural numbers and $0 \leq b_{n+1}<b_{n}$, there exists $N \in \mathbb{N}$ such that $b_{N}=0$. It proves that the algorithm ends after finitely many steps. Furthermore

$$
\operatorname{gcd}(a, b)=\operatorname{gcd}\left(a_{0}, b_{0}\right)=\operatorname{gcd}\left(a_{1}, b_{1}\right)=\cdots=\operatorname{gcd}\left(a_{N}, b_{N}\right)=a_{N}
$$

So the algorithm computes $\operatorname{gcd}(a, b)$ as expected.
Algorithm: Euclid's algorithm in pseudocode
Result: $\operatorname{gcd}(a, b)$ where $a, b \in \mathbb{Z}$ not both zero.
$a \leftarrow|a|$
$b \leftarrow|b|$
while $b \neq 0$ do
$r \leftarrow a \% b$ ( the remainder of the Euclidean division $a=b q+r$ with $0 \leq r<b$ )
$a \leftarrow b$
$b \leftarrow r$
end
return $a$

Example 36. We want to compute $\operatorname{gcd}(600,-136)$ :


Hence $\operatorname{gcd}(600,-136)=8$.
It is possible to obtain a suitable Bézout's identity from the above algorithm by going backward.

$$
\begin{aligned}
8 & =56+24 \times(-2) \\
& =56+(136+56 \times(-2)) \times(-2) \\
& =136 \times(-2)+56 \times 5 \\
& =136 \times(-2)+(600+136 \times(-4 \\
8 & =600 \times 5+(-136) \times 22
\end{aligned}
$$

$$
=136 \times(-2)+(600+136 \times(-4)) \times 5 \quad \text { since } 56=600-136 \times 4
$$

## 7 Coprime integers

Definition 37. Let $a, b \in \mathbb{Z}$ not both zero. We say that $a$ and $b$ are coprime (or relatively prime) if $\operatorname{gcd}(a, b)=1$.
The following result states that the converse of Bézout's identity holds for coprime numbers.
Proposition 38. Let $a, b \in \mathbb{Z}$ not both zero. Then

$$
\operatorname{gcd}(a, b)=1 \Leftrightarrow \exists u, v \in \mathbb{Z}, a u+b v=1
$$

Proof.
$\Rightarrow$ : it is simply Bézout's identity.
$\Leftarrow$ : let $a, b \in \mathbb{Z}$ not both zero. Assume that $a u+b v=1$ for some $u, v \in \mathbb{Z}$.
Set $d=\operatorname{gcd}(a, b)$. Then $d \mid a$ and $d \mid b$, hence $d \mid(a u+b v)=1$. So $|d|=1$. But since $d \in \mathbb{N}$, we get that $d=1$.
Theorem 39 (Gauss' lemma). $\forall a, b, c \in \mathbb{Z},\left\{\left.\begin{array}{l}\operatorname{gcd}(a, b)=1 \\ a \mid b c\end{array} \Longrightarrow a \right\rvert\, c\right.$
Proof. Let $a, b, c \in \mathbb{Z}$ such that $\operatorname{gcd}(a, b)=1$ and $a \mid b c$. Then there exists $k \in \mathbb{Z}$ such that $b c=k a$. By Bézout's identity, there exist $u, v \in \mathbb{Z}$ such that $1=a u+b v$.
Thus $c=(a u+b v) c=a u c+b c v=a u c+k a v=a(u c+k v)$. Hence $a \mid c$.
The following result is very useful.
Proposition 40. Let $a, b, c \in \mathbb{Z}$. If $a|c, b| c$ and $\operatorname{gcd}(a, b)=1$ then $a b \mid c$.
Proof. Since $a \mid c$ and $b \mid c$, there exist $k, l \in \mathbb{Z}$ such that $c=a k$ and $c=b l$.
Since $\operatorname{gcd}(a, b)=1$, by Bézout's identity, there exists $u, v \in \mathbb{Z}$ such that $a u+b v=1$.
Then $c=a u c+b v c=a u b l+b v a k=a b(u l+v k)$, so that $a b \mid c$.

## 8 A diophantine equation

Theorem 41. Let $a, b, c \in \mathbb{Z}$ with $a$ and $b$ not both zero.
Then the equation $a x+b y=c$ has an integer solution if and only if $\operatorname{gcd}(a, b) \mid c$.
Proof.
$\Rightarrow$ : Assume that $a x+b y=c$ for some $(x, y) \in \mathbb{Z}^{2}$.
Since $\operatorname{gcd}(a, b) \mid a$ and $\operatorname{gcd}(a, b) \mid b$, we get that $\operatorname{gcd}(a, b) \mid a x+b y=c$.
$\Leftarrow$ : Assume that $\operatorname{gcd}(a, b) \mid c$, then there exists $k \in \mathbb{Z}$ such that $c=k \operatorname{gcd}(a, b)$.
By Bézout's identity, there exists $(u, v) \in \mathbb{Z}^{2}$ such that $a u+b v=\operatorname{gcd}(a, b)$ hence $a k u+b k v=k \operatorname{gcd}(a, b)=c$. Therefore $(k u, k v)$ is an integer solution of the equation.

How to find all the integer solutions of an equation of the form $a x+b y=c$ with $a \neq 0, b \neq 0$ and $\operatorname{gcd}(a, b) \mid c$ ?

- Step 1: reduction to the case where $\operatorname{gcd}(a, b)=1$.

There exist $\tilde{a}, \tilde{b}, \tilde{c} \in \mathbb{Z}$ such that $a=\tilde{a} \operatorname{gcd}(a, b), b=\tilde{b} \operatorname{gcd}(a, b)$ and $c=\tilde{c} \operatorname{gcd}(a, b)$.
Hence $a x+b y=c \Leftrightarrow \tilde{a} x+\tilde{b} y=\tilde{c}$.
Note that $\operatorname{gcd}(a, b)=\operatorname{gcd}(\tilde{a} \operatorname{gcd}(a, b), \tilde{b} \operatorname{gcd}(a, b))=\operatorname{gcd}(a, b) \operatorname{gcd}(\tilde{a}, \tilde{b})$. Hence $\operatorname{gcd}(\tilde{a}, \tilde{b})=1$.

- Step 2: find a first solution.

By Bézout's identity, there exist $u, v \in \mathbb{Z}$ such that $\tilde{a} u+\tilde{b} v=1$ (we may find such a couple $(u, v)$ using Euclid's algorithm).
Thence $\tilde{a} \tilde{c} u+\tilde{b} \tilde{c} v=\tilde{c}$. Therefore we obtain a solution $\left(x_{0}, y_{0}\right)=(\tilde{c} u, \tilde{c} v)$ of $\tilde{a} x+\tilde{b} y=\tilde{c}$.

- Step 3: study the other solutions.

Let $(x, y) \in \mathbb{Z}^{2}$ satisfying $\tilde{a} x+\tilde{b} y=\tilde{c}$. Then $\tilde{a}\left(x-x_{0}\right)+\tilde{b}\left(y-y_{0}\right)=0$, i.e. $\tilde{b}\left(y-y_{0}\right)=\tilde{a}\left(x_{0}-x\right)$.
Since $\tilde{a} \mid \tilde{b}\left(y-y_{0}\right)$ and $\operatorname{gcd}(\tilde{a}, \tilde{b})=1$, by Gauss' lemma, $\tilde{a} \mid y-y_{0}$, i.e. there exists $k \in \mathbb{Z}$ such that $k \tilde{a}=y-y_{0}$, i.e. $y=y_{0}+k \tilde{a}$.

Then $\tilde{a}\left(x_{0}-x\right)=\tilde{b}\left(y-y_{0}\right)=k \tilde{a} \tilde{b}$. Since $\tilde{a} \neq 0$, we get $x_{0}-x=k \tilde{b}$, i.e. $x=x_{0}-k \tilde{b}$.
We proved that there exists $k \in \mathbb{Z}$ such that $(x, y)=\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right)$.

- Step 4: check the converse!

We proved that if $(x, y) \in \mathbb{Z}^{2}$ is a solution, then there exists $k \in \mathbb{Z}$ such that $(x, y)=\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right)$. It means that the solutions are among $(x, y) \in\left\{\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right): k \in \mathbb{Z}\right\}$.
Otherwise stated, it means that $\left\{(x, y) \in \mathbb{Z}^{2}: \tilde{a} x+\tilde{b} y=\tilde{c}\right\} \subset\left\{\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right): k \in \mathbb{Z}\right\}$.
It doesn't mean that they are all solutions, we need to check that separately, i.e. we need to prove the other inclusion.
Conversely, let's prove that for every $k \in \mathbb{Z},(x, y)=\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right)$ is a solution:

$$
\tilde{a}\left(x_{0}-k \tilde{b}\right)+\tilde{b}\left(y_{0}+k \tilde{a}\right)=\tilde{a} x_{0}+\tilde{b} y_{0}=\tilde{c}
$$

- Step 5: Conclusion!

The solutions are exactly the $(x, y)=\left(x_{0}-k \tilde{b}, y_{0}+k \tilde{a}\right)$ for $k \in \mathbb{Z}$.

See Slide 6 of Lecture 7 (Feb 2) for a concrete example.

## A Appendix: properties of the strict order

Recall that given $a, b \in \mathbb{Z}, a<b$ means ( $a \leq b$ and $a \neq b$ ).
The following properties of $<$ are easy to derive from the ones of $\leq$.

- $\forall a, b, c \in \mathbb{Z},(a<b$ and $b \leq c) \Longrightarrow a<c$
- $\forall a, b, c \in \mathbb{Z},(a \leq b$ and $b<c) \Longrightarrow a<c$
- $\forall a, b, c, d \in \mathbb{Z},(a<b$ and $c \leq d) \Longrightarrow a+c<b+d$
- $\forall a, b, c \in \mathbb{Z}, a<b \Longrightarrow a+c<b+c$
(that's a special case of the previous one where $d=c$ )
- $\forall a, b, c \in \mathbb{Z},(a<b$ and $c>0) \Longrightarrow a c<b c$
- $\forall a, b, c \in \mathbb{Z},(a<b$ and $c<0) \Longrightarrow a c>b c$
- $\forall a, b \in \mathbb{Z}, a<b \Leftrightarrow a+1 \leq b$
- Given $a, b \in \mathbb{Z}$, exactly one of the following occurs:
(i) $a<b$
(ii) $a=b$
(iii) $a>b$

Particularly, the negation of $a \leq b$ if $a>b$.

## B Appendix: implementation of Euclid's algorithm in Julia

## Euclid's algorithm in Julia (iterative)

```
function euclid(a::Integer, b::Integer)
    a != 0 || b != 0 || error("a and b must not be both zero")
    a = abs(a)
    b = abs (b)
    while b != 0
        r = a%b
        a = b
        b}=
    end
    return a
end
```

Actually, it is not important to replace $a$ and $b$ by their respective absolute values in the initialization. In this case, the sequence $\left(b_{n}\right)$ is eventually non-negative so the algorithm stops as earlier and we just have to make sure that we return the absolute value of $a$ at the end.

That being said, you should be careful because most programming languages don't use the above convention for Euclidean division. Instead, they require the remainder $r$ to have the same sign as $b$, i.e. $r$ satisfies $0 \leq r<b$ if $b>0$ or $b<r \leq 0$ if $b<0$.

But it doesn't matter for Euclid's algorithm: indeed, with this convention, the sequence $\left|b_{n}\right|$ is still decreasing, so the algorithm stops.

Therefore, we can simply write the following program (here $\operatorname{gcd}(0,0)=0$ by convention).
Euclid's algorithm in Julia (recursive)

```
function euclid(a::Integer, b::Integer)
    b != 0 || return abs(a)
    return euclid(b, a%b)
end
```


[^0]:    ${ }^{1}$ When $b>0$, the informal idea of this proof consists in determining how many times we can add before exceeding $a$, which will give the quotient. Then the remainder will be obtained by filling the difference in order to reach $a$.
    Intuitively, if the quotient exists, it has to be the greatest $p$ such that $b p \leq a$. We have to prove the existence of such a number and then to check formally that this idea is actually correct.

